JavaScript Static Analysis for Evolving Language Specifications

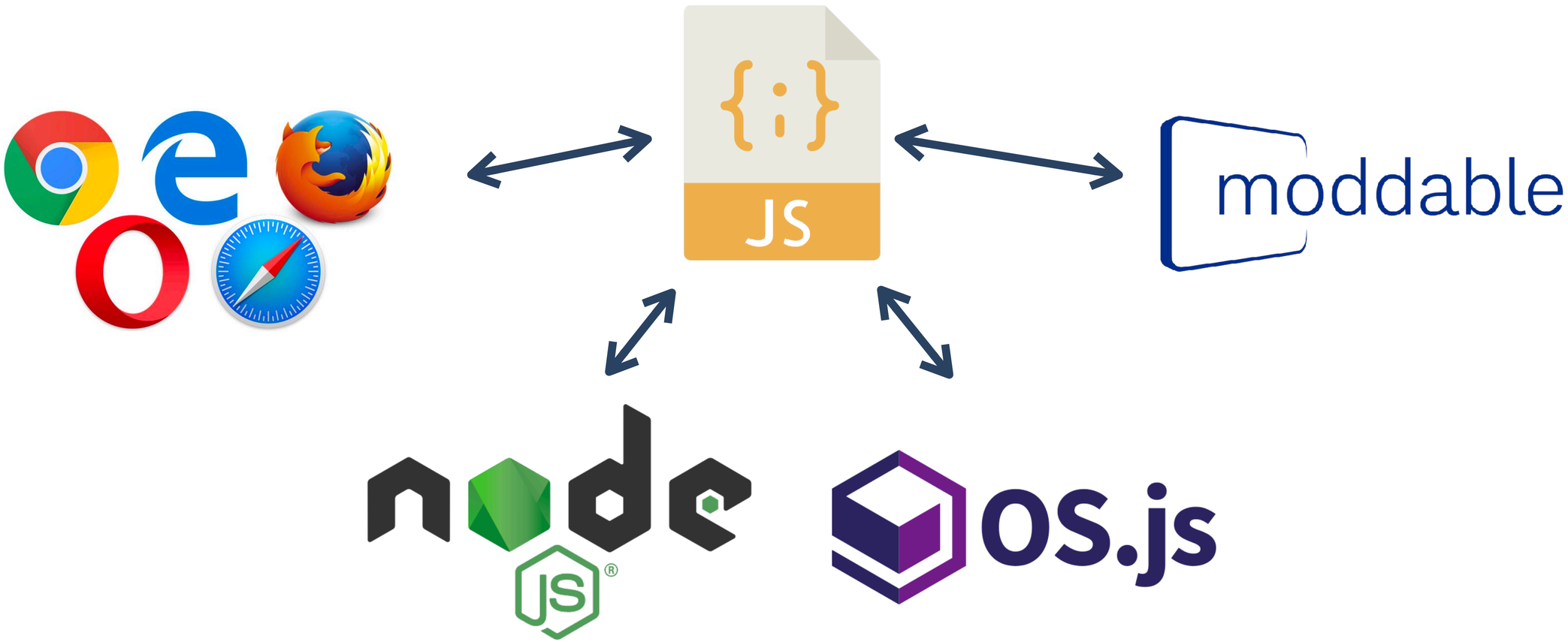
[SW재난연구센터] 겨울정기워크샵

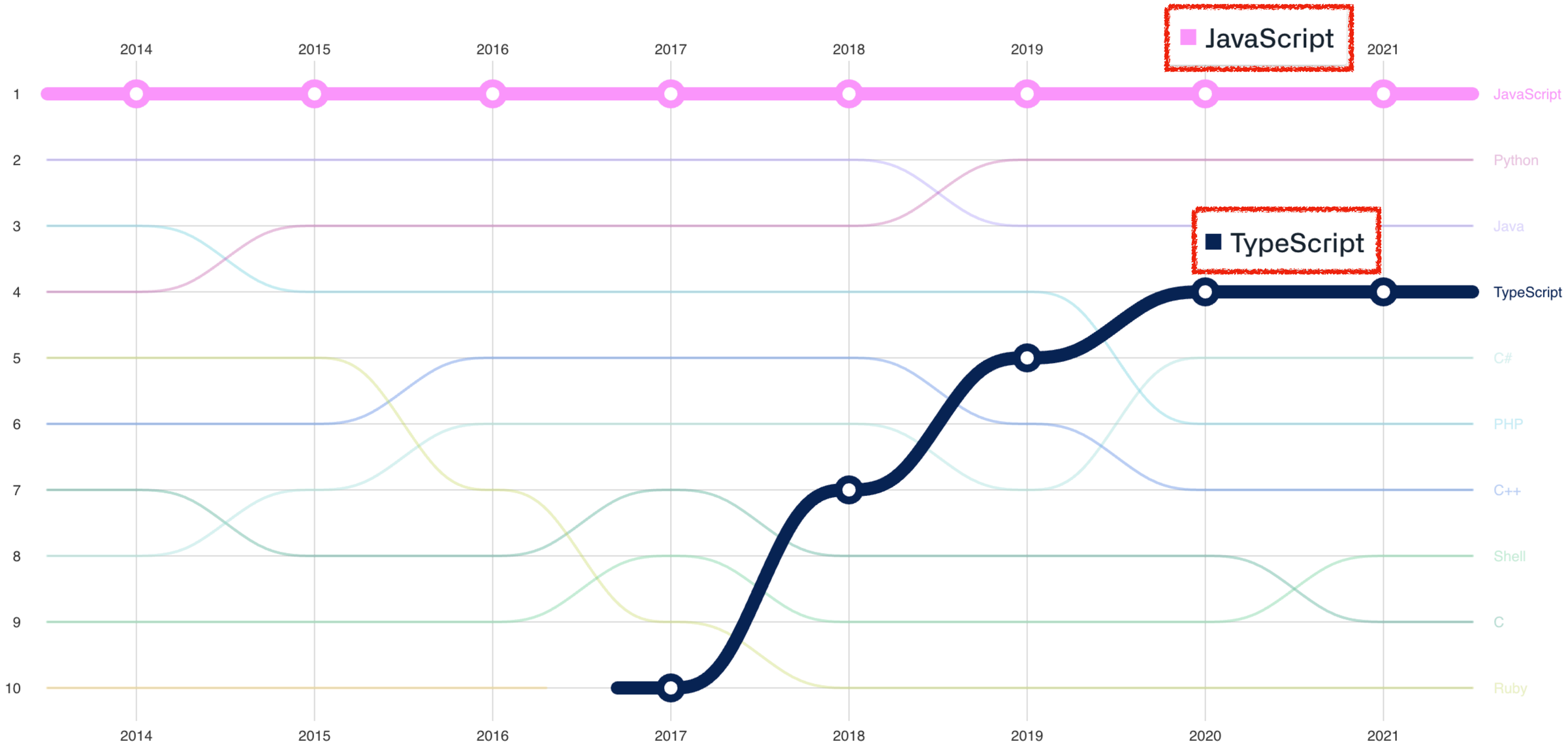
Jihyeok Park

PLRG @ KAIST

February 9, 2022

JavaScript Is Everywhere





https://octoverse.github.com/

JavaScript Complex Semantics

```
function f(x) { return x == !x; }
```

Always return `false`?

JavaScript Complex Semantics

```
function f(x) { return x == !x; }
```

Always return `false`?

NO!!

```
f([]) -> [] == ![]
      -> [] == false
      -> +[] == +false
      -> 0 == 0
      -> true
```

ECMA-262: ECMAScript Specification

Syntax

ArrayLiteral[Yield, Await] :

[*Elision*opt]

[*ElementList*[?Yield, ?Await]]

[*ElementList*[?Yield, ?Await] , *Elision*opt]

The production of *ArrayLiteral* in ES12

Semantics

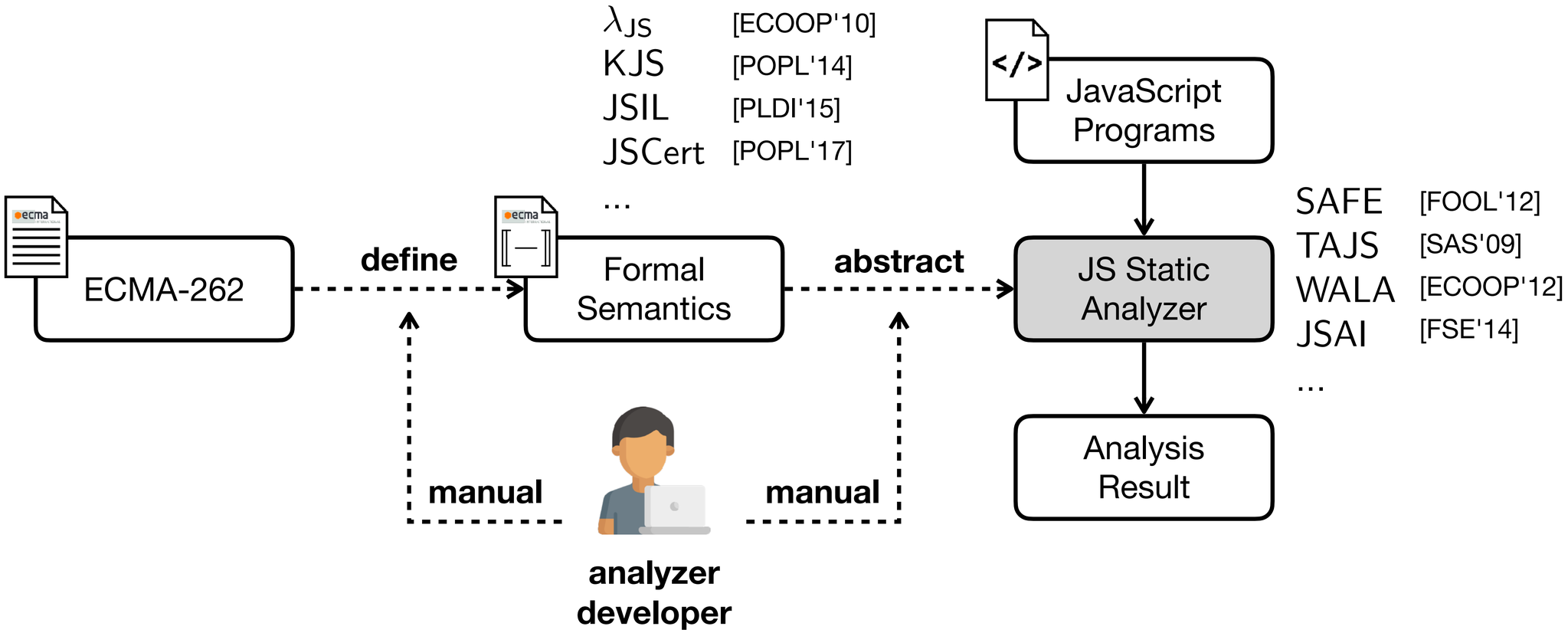
13.2.5.2 Runtime Semantics: Evaluation

ArrayLiteral : [*ElementList* , *Elision*opt]

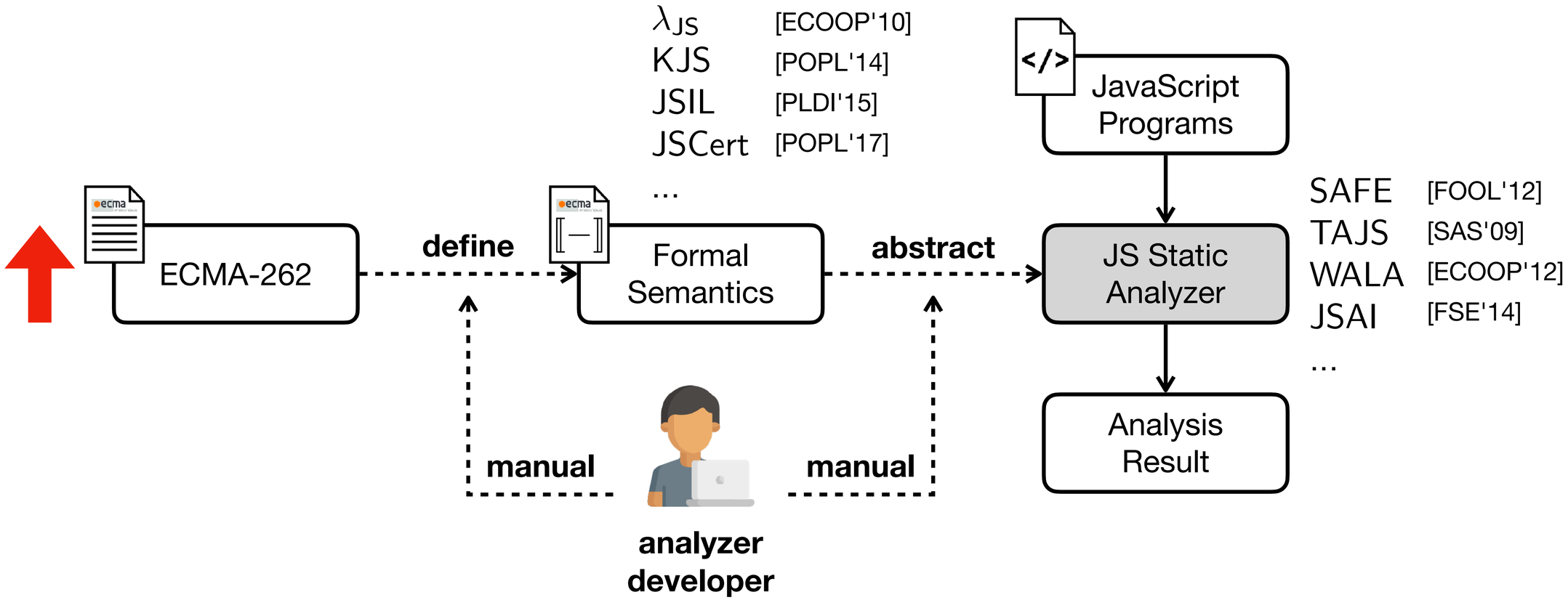
1. Let *array* be ! ArrayCreate(0).
2. Let *nextIndex* be the result of performing ArrayAccumulation for *ElementList* with arguments *array* and 0.
3. ReturnIfAbrupt(*nextIndex*).
4. If *Elision* is present, then
   a. Let *len* be the result of performing ArrayAccumulation for *Elision* with arguments *array* and *nextIndex*.
   b. ReturnIfAbrupt(*len*).
5. Return *array*.

The Evaluation algorithm for the third alternative of *ArrayLiteral* in ES12

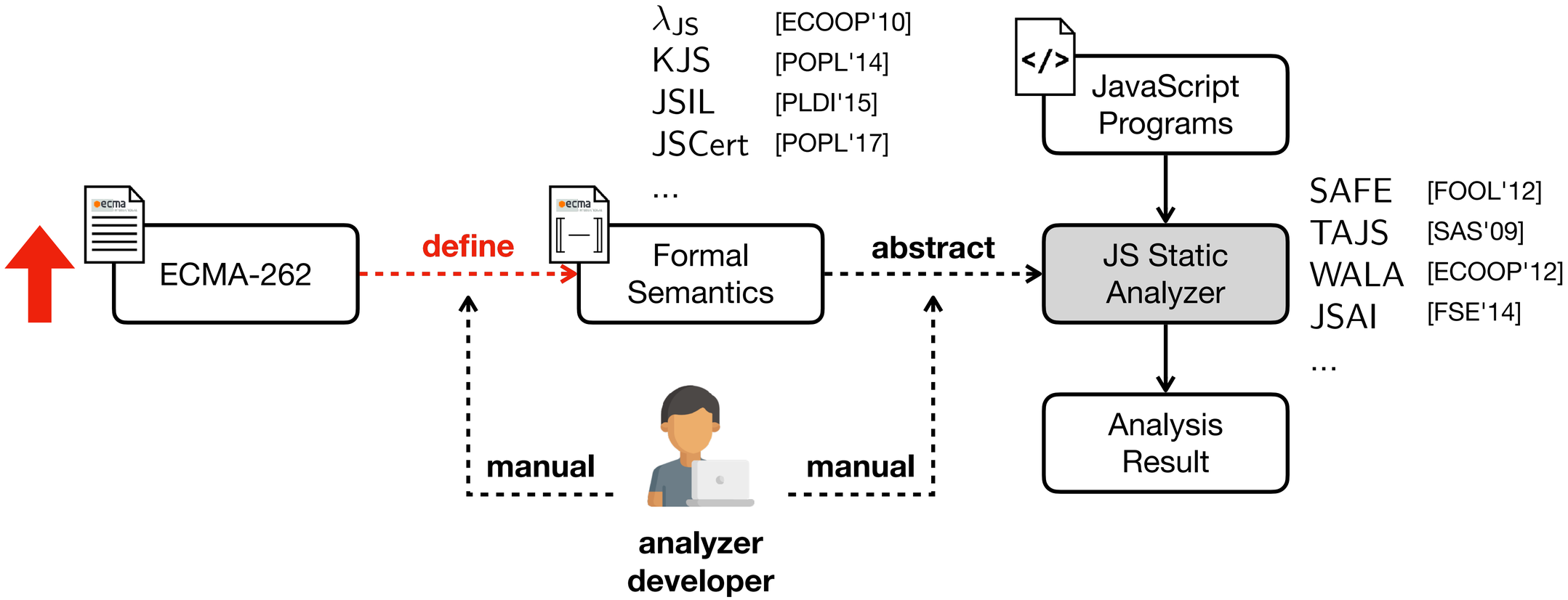
Problem: Manual JavaScript Static Analyzer



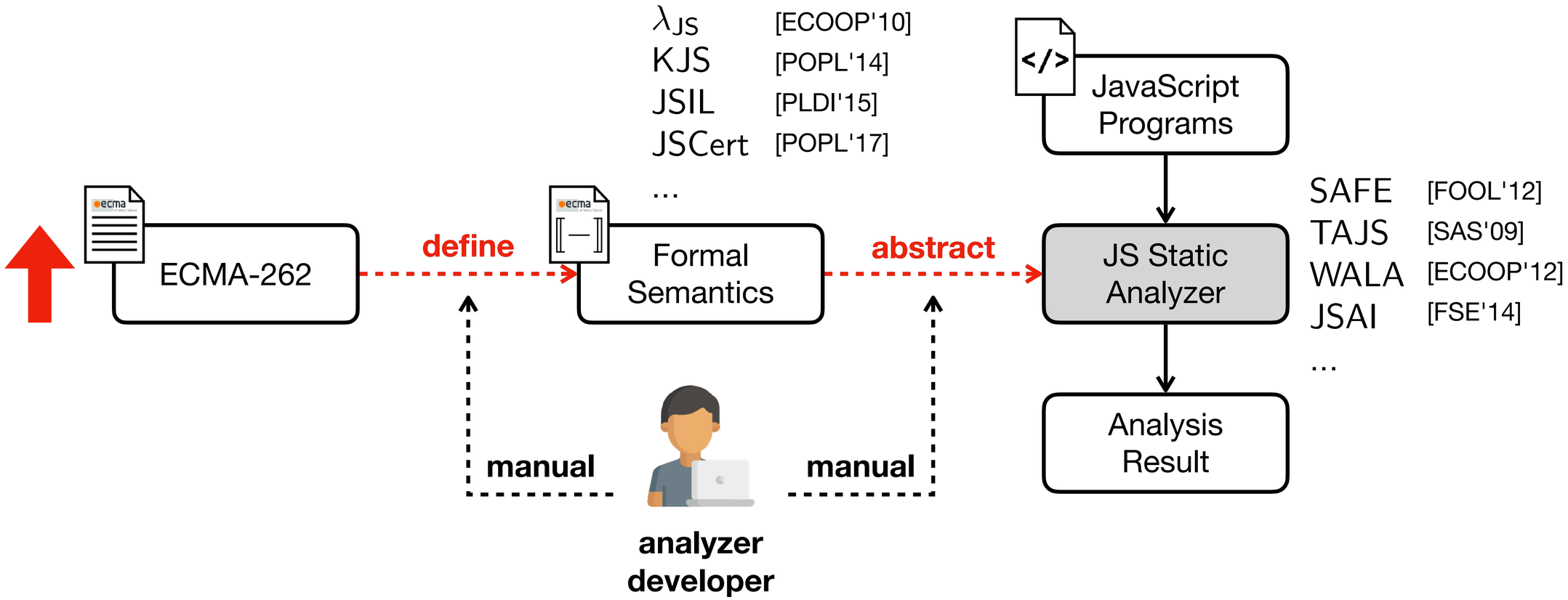
Problem: Manual JavaScript Static Analyzer



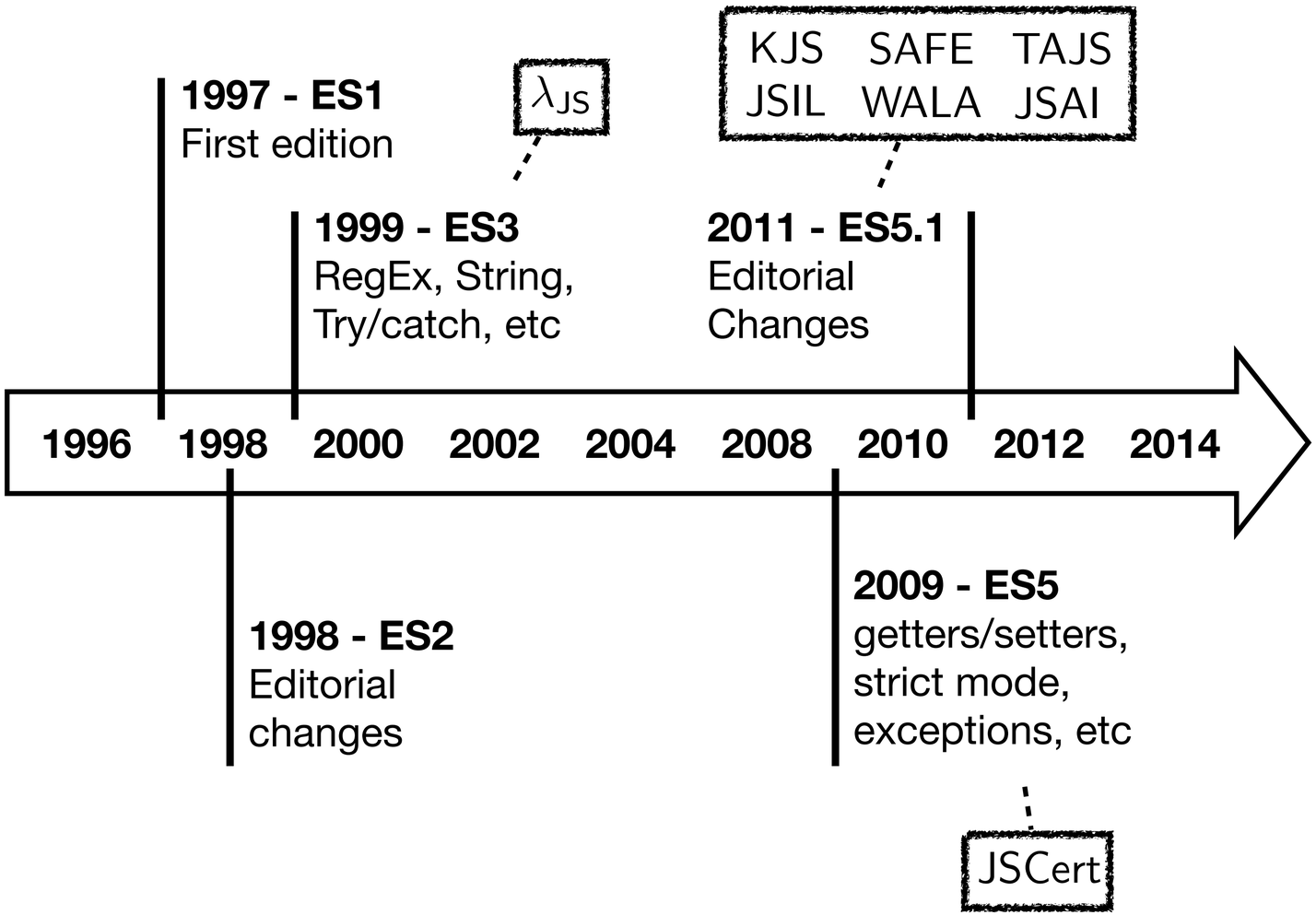
Problem: Manual JavaScript Static Analyzer



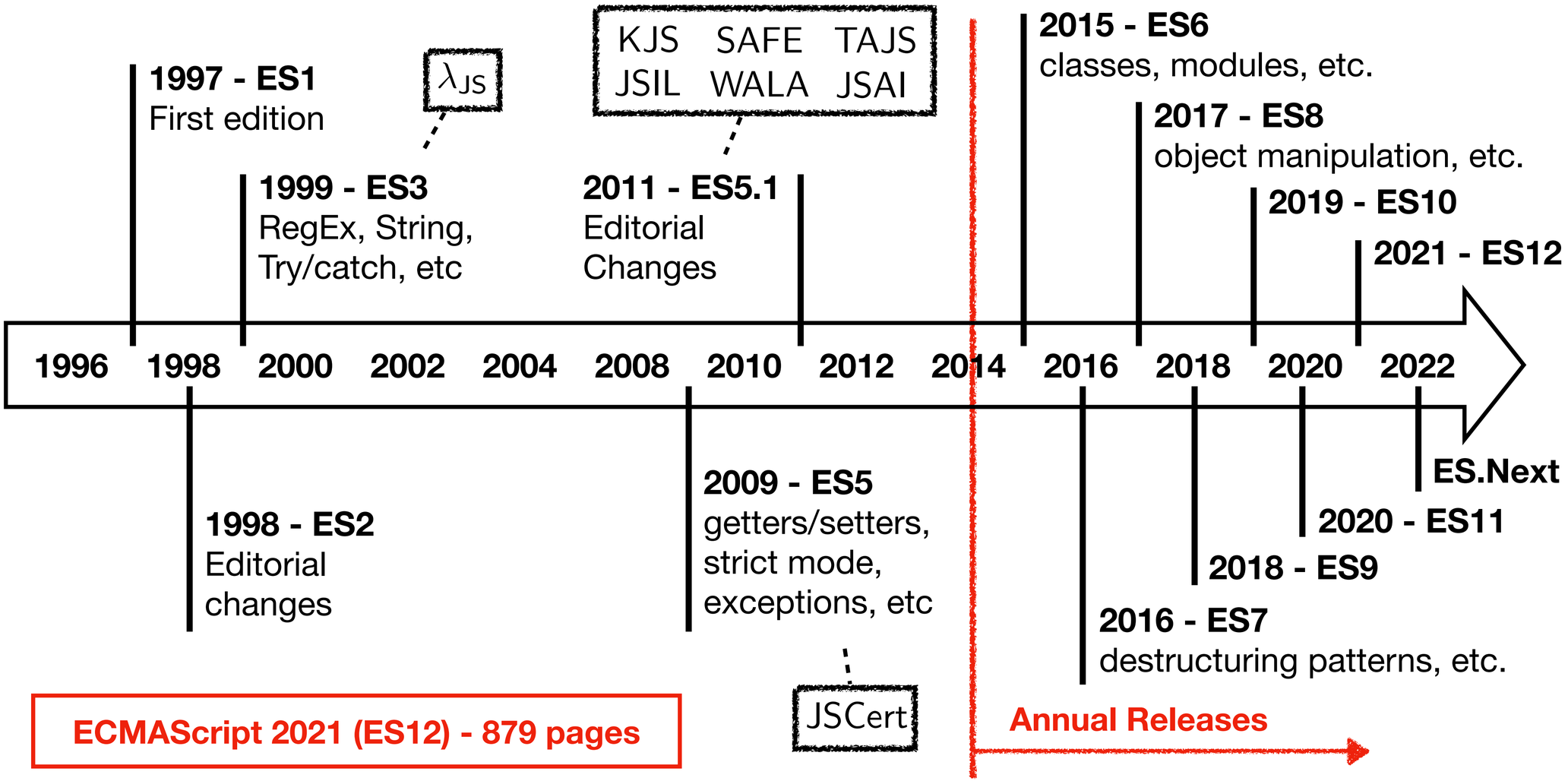
Problem: Manual JavaScript Static Analyzer



Problem: Fast Evolving JavaScript

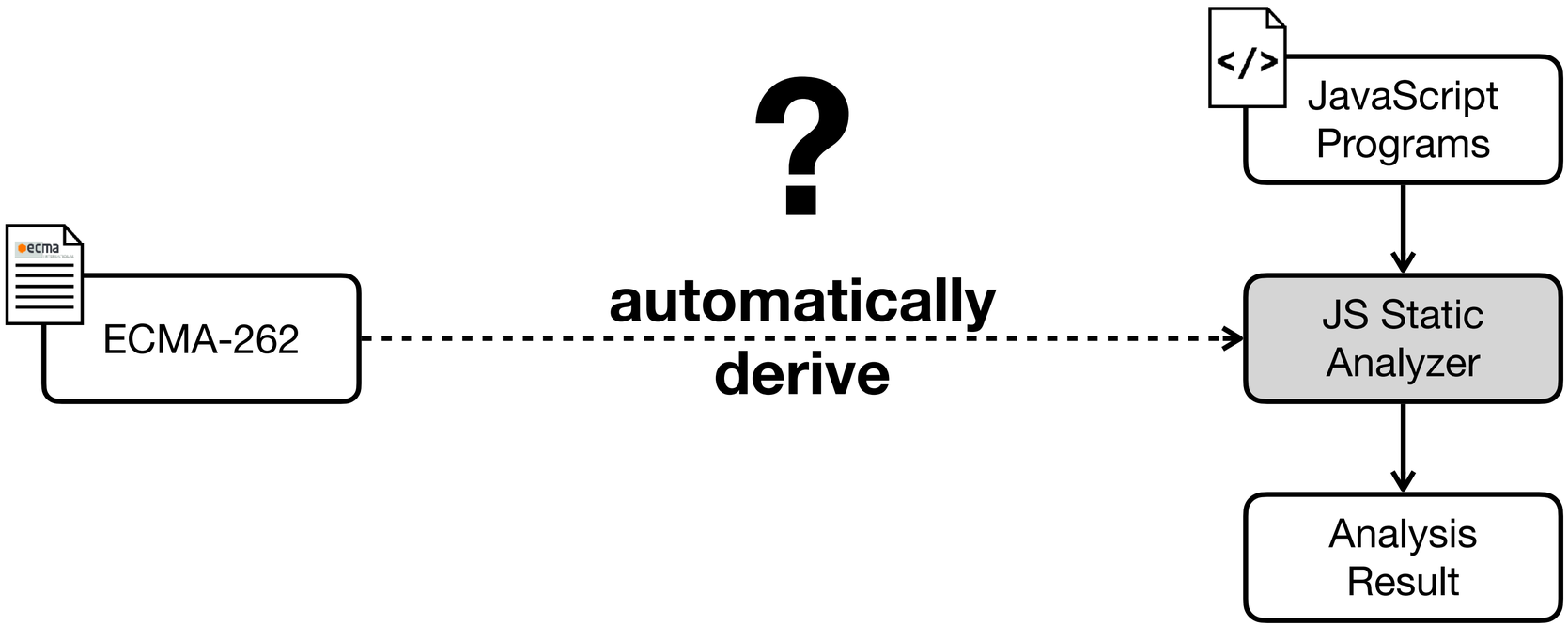


Problem: Fast Evolving JavaScript

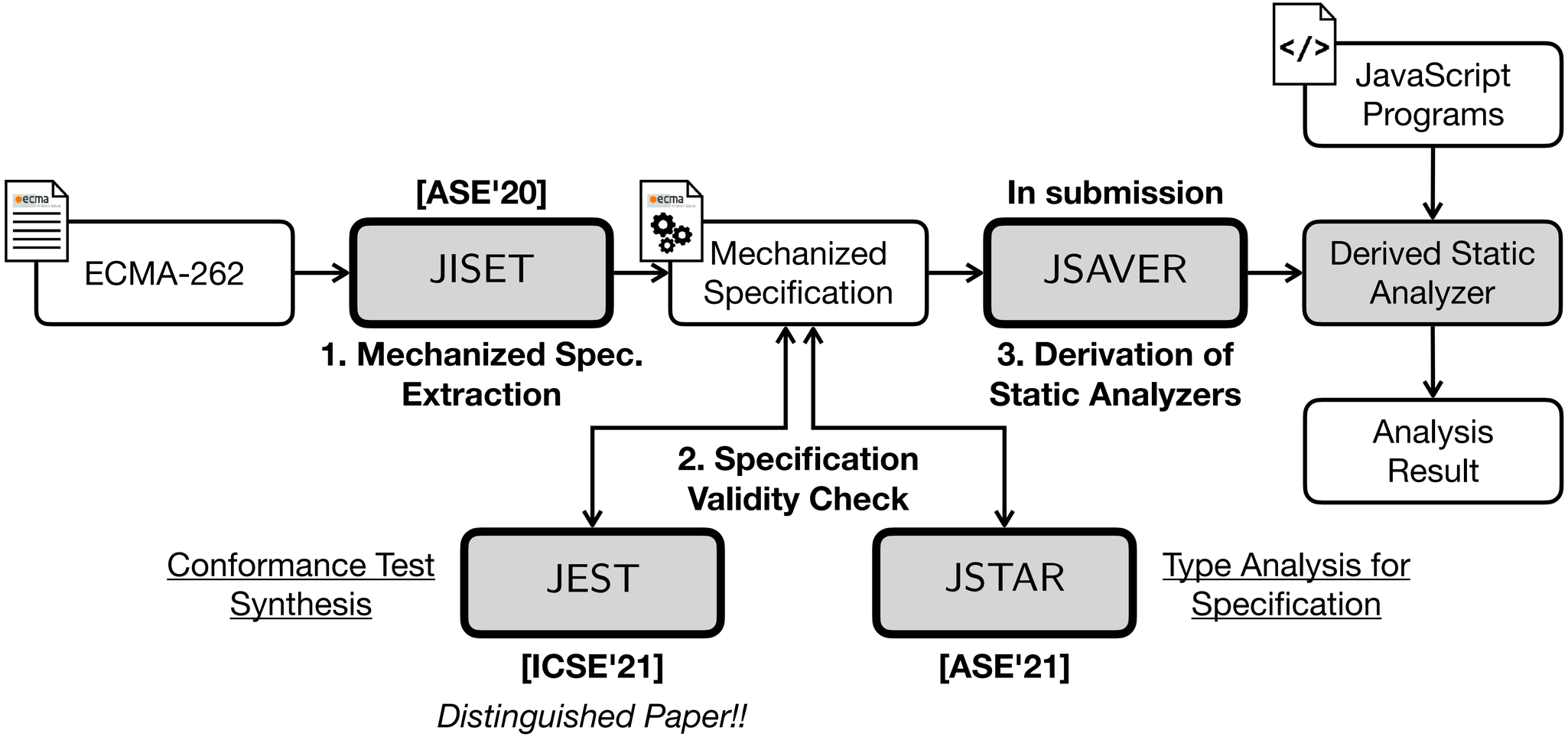


ECMAScript 2021 (ES12) - 879 pages

Main Idea: Deriving Static Analyzer from Spec.



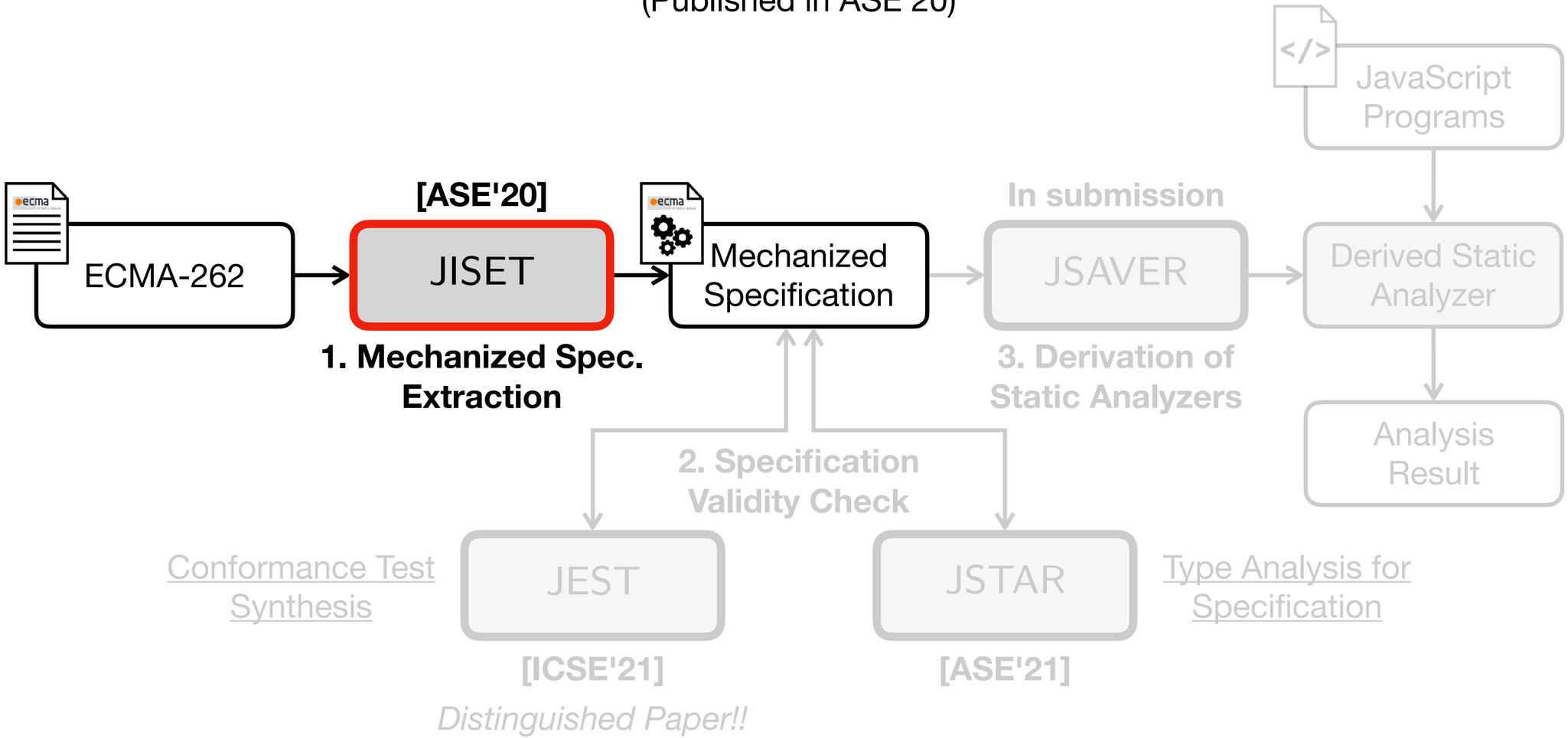
Overall Structure



JISET: JavaScript IR-based Semantics Extraction Toolchain

Jihyeok Park, Jihee Park, Seungmin An, and Sukyoung Ryu

(Published in ASE'20)



Motivation: Patterns in Writing Style of ECMA-262

13.2.5.2 Runtime Semantics: Evaluation

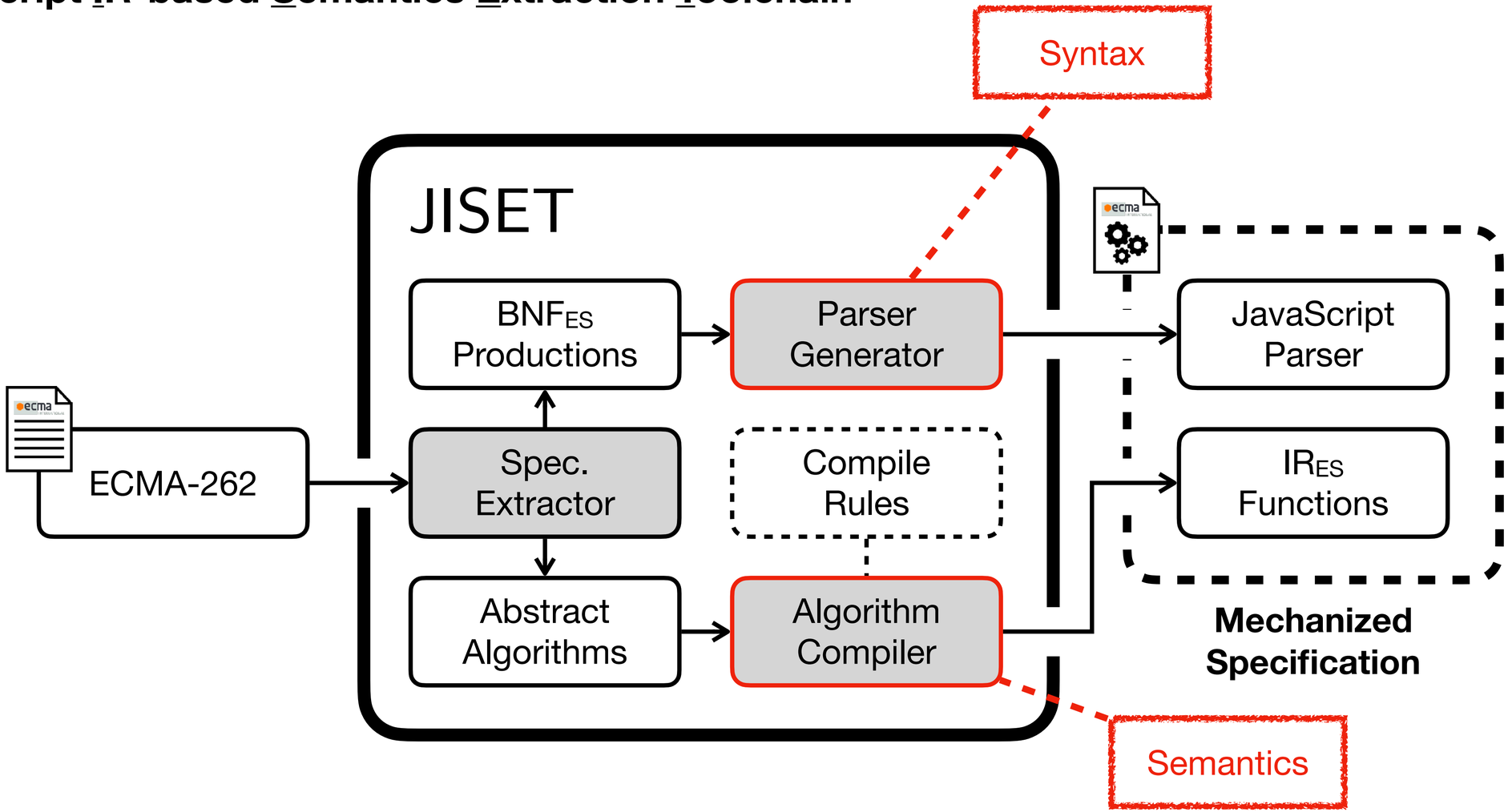
ArrayLiteral : [*ElementList* , *Elision*$_{opt}$]

1. Let *array* be ! ArrayCreate(0).
2. Let *nextIndex* be the result of performing ArrayAccumulation for *ElementList* with arguments *array* and 0.
3. ReturnIfAbrupt(*nextIndex*).
4. If *Elision* is present, then
   a. Let *len* be the result of performing ArrayAccumulation for *Elision* with arguments *array* and *nextIndex*.
   b. ReturnIfAbrupt(*len*).
5. Return *array*.

The Evaluation **algorithm for the third alternative of** *ArrayLiteral* **in ES12**

JISET [ASE'20]

JavaScript IR-based Semantics Extraction Toolchain



JISET - Parser Generator (Syntax)

ArrayLiteral[Yield, Await] :

[*Elision*$_{opt}$]

[*ElementList*[?Yield, ?Await]]

[*ElementList*[?Yield, ?Await] , *Elision*$_{opt}$]

JavaScript Parser in Scala

Parsing Expression Grammar
(+ Lookahead Parsing)

```
val ArrayLiteral: List[Boolean] => LAParser[T] = memo {
  case List(Yield, Await) =>
  "[" ~ opt(Elision) ~ "]"                ^^ ArrayLiteral0 |
  "[" ~ ElementList(Yield,Await) ~ "]"  ^^ ArrayLiteral1 |
  "[" ~ ElementList(Yield,Await) ~ ","
      ~ opt(Elision)~ "]"                 ^^ ArrayLiteral2
}
```

(POPL'04) Bryan Ford, "Parsing Expression Grammars: A Recognition-based Syntactic Foundation"

JISET - Algorithm Compiler (Semantics)

13.2.5.2 Runtime Semantics: Evaluation

ArrayLiteral : [*ElementList* , *Elision*$_{opt}$]

1. Let *array* be ! ArrayCreate(0).
2. Let *nextIndex* be the result of performing ArrayAccumulation for *ElementList* with arguments *array* and 0.
3. ReturnIfAbrupt(*nextIndex*).
4. If *Elision* is present, then
   a. Let *len* be the result of performing ArrayAccumulation for *Elision* with arguments *array* and *nextIndex*.
   b. ReturnIfAbrupt(*len*).
5. Return *array*.

IR$_{ES}$ Functions

118 Compile Rules for Steps in Abstract Algorithms

```
syntax def ArrayLiteral[2].Evaluation(
  this, ElementList, Elision
) {
  let array = [! (ArrayCreate 0)]
  let nextIndex = (ElementList.ArrayAccumulation array 0)
  [? nextIndex]
  if (! (= Elision absent)) {
    let len = (Elision.ArrayAccumulation array nextIndex)
    [? len]
  }
  return array
}
```

JISET - Evaluation

≈ 95% Compiled

auto manual

T: Total **L**: Core Language Semantics **B**: Built-in Libraries

Version	# Algo.		
ES7	2,105	T	**10,471 / 10,982** (95.35%)
		L	**8,041 / 8,415** (95.56%)
		B	**2,430 / 2,567** (94.66%)
ES8	2,238	T	**11,181 / 11,732** (95.30%)
		L	**8,453 / 8,811** (95.94%)
		B	**2,728 / 2,921** (93.39%)
ES9	2,370	T	**11,849 / 12,393** (95.61%)
		L	**8,932 / 9,311** (95.93%)
		B	**2,917 / 3,082** (94.65%)
ES10	2,396	T	**12,022 / 12,569** (95.65%)
		L	**9,073 / 9,456** (94.95%)
		B	**2,949 / 3,113** (94.73%)
ES11	2,521	T	**12,505 / 13,047** (94.85%)
		L	**9,495 / 9,881** (96.09%)
		B	**3,010 / 3,166** (95.07%)
ES12	2,640	T	**12,975 / 13,544** (95.80%)
		L	**9,717 / 10,136** (95.87%)
		B	**3,258 / 3,408** (95.60%)
Average	2,378	T	**11,834 / 12,378** (95.61%)
		L	**8,952 / 9,335** (95.90%)
		B	**2,882 / 3,043** (94.71%)

JISET - Evaluation

≈ 95% Compiled

auto manual

T: Total L: Core Language Semantics B: Built-in Libraries

Version	# Algo.		
ES7	2,105	T	**10,471 / 10,982** (95.35%)
		L	**8,041 / 8,415** (95.56%)
		B	**2,430 / 2,567** (94.66%)
ES8	2,238	T	**11,181 / 11,732** (95.30%)
		L	**8,453 / 8,811** (95.94%)
		B	**2,728 / 2,921** (93.39%)
ES9	2,370	T	**11,849 / 12,393** (95.61%)
		L	**8,932 / 9,311** (95.93%)
		B	**2,917 / 3,082** (94.65%)
ES10	2,396	T	**12,022 / 12,569** (95.65%)
		L	**9,073 / 9,456** (94.95%)
		B	**2,949 / 3,113** (94.73%)
ES11	2,521	T	**12,505 / 13,047** (94.85%)
		L	**9,495 / 9,881** (96.09%)
		B	**3,010 / 3,166** (95.07%)
ES12	2,640	T	**12,975 / 13,544** (95.80%)
		L	**9,717 / 10,136** (95.87%)
		B	**3,258 / 3,408** (95.60%)
Average	2,378	T	**11,834 / 12,378** (95.61%)
		L	**8,952 / 9,335** (95.90%)
		B	**2,882 / 3,043** (94.71%)

Complete Missing Parts

JISET - Evaluation

≈ 95% Compiled

Passed All Tests

auto manual

T: Total L: Core Language Semantics B: Built-in Libraries

Version	# Algo.		
ES7	2,105	T	**10,471 / 10,982** (95.35%)
		L	**8,041 / 8,415** (95.56%)
		B	**2,430 / 2,567** (94.66%)
ES8	2,238	T	**11,181 / 11,732** (95.30%)
		L	**8,453 / 8,811** (95.94%)
		B	**2,728 / 2,921** (93.39%)
ES9	2,370	T	**11,849 / 12,393** (95.61%)
		L	**8,932 / 9,311** (95.93%)
		B	**2,917 / 3,082** (94.65%)
ES10	2,396	T	**12,022 / 12,569** (95.65%)
		L	**9,073 / 9,456** (94.95%)
		B	**2,949 / 3,113** (94.73%)
ES11	2,521	T	**12,505 / 13,047** (94.85%)
		L	**9,495 / 9,881** (96.09%)
		B	**3,010 / 3,166** (95.07%)
ES12	2,640	T	**12,975 / 13,544** (95.80%)
		L	**9,717 / 10,136** (95.87%)
		B	**3,258 / 3,408** (95.60%)
Average	2,378	T	**11,834 / 12,378** (95.61%)
		L	**8,952 / 9,335** (95.90%)
		B	**2,882 / 3,043** (94.71%)

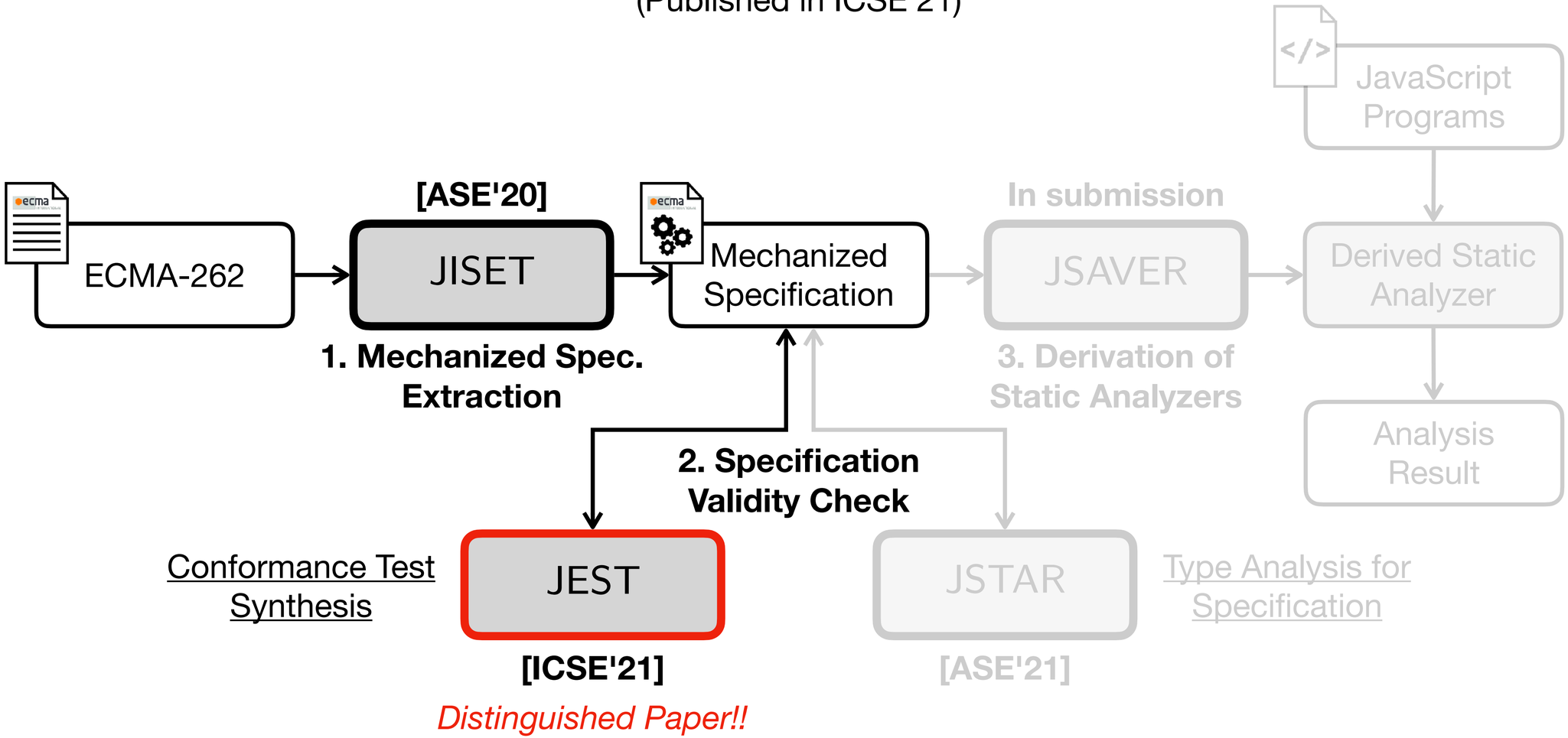
Complete Missing Parts

- **Test262**
  (Official Conformance Tests)
  - 18,556 applicable tests
- **Parsing tests**
  - Passed all 18,556 tests
- **Evaluation Tests**
  - Passed all 18,556 tests

JEST: N+1-version Differential Testing of Both JavaScript Engines

Jihyeok Park, Seungmin An, Dongjun Youn, Gyeongwon Kim, and Sukyoung Ryu

(Published in ICSE'21)



JEST - Conformance with Engines

ECMA-262

?

Conform

GraalVM

QuickJS

moddable

JavaScript Engines

JEST - N+1-version Differential Testing

ECMA-262

Synthesize

Test

test

test

test

test

GraalVM

QuickJS

moddable

JavaScript Engines

JEST - N+1-version Differential Testing

ECMA-262

Synthesize

Test

test

test

test

test

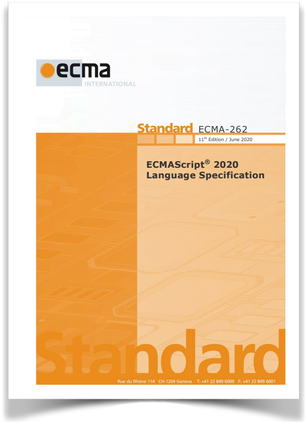
GraalVM

QuickJS

moddable

JavaScript Engines

JEST - N+1-version Differential Testing



ECMA-262

Synthesize → Test → test, test, test, test

GraalVM

QuickJS

moddable

JavaScript Engines

An engine bug in

JEST - N+1-version Differential Testing

ECMA-262

Synthesize

Test

test

test

test

test

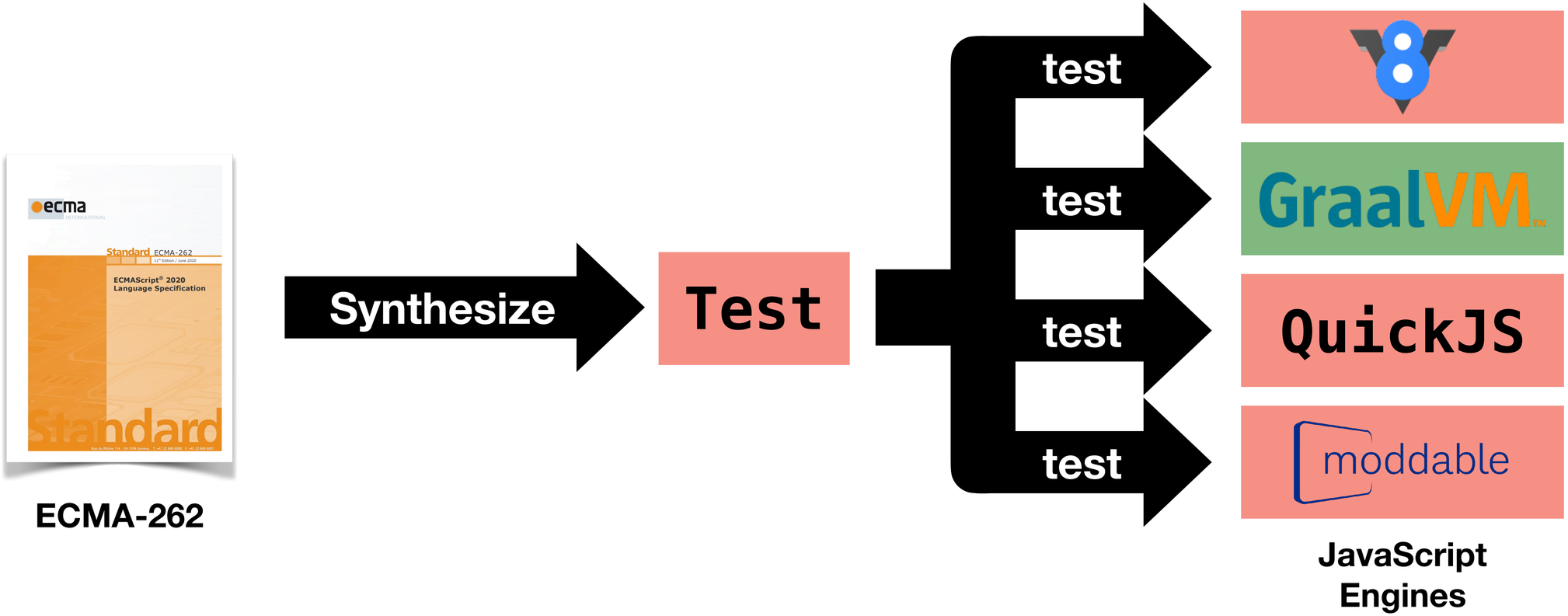
GraalVM

QuickJS

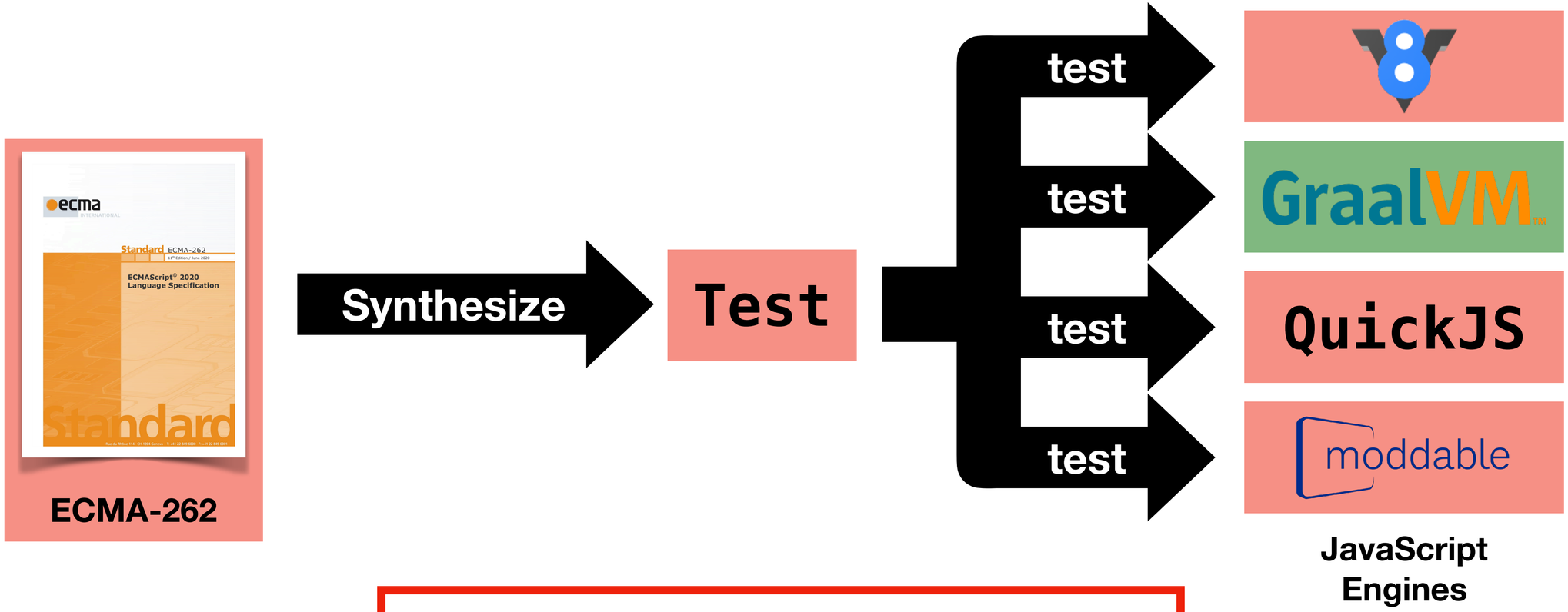
moddable

JavaScript Engines

JEST - N+1-version Differential Testing

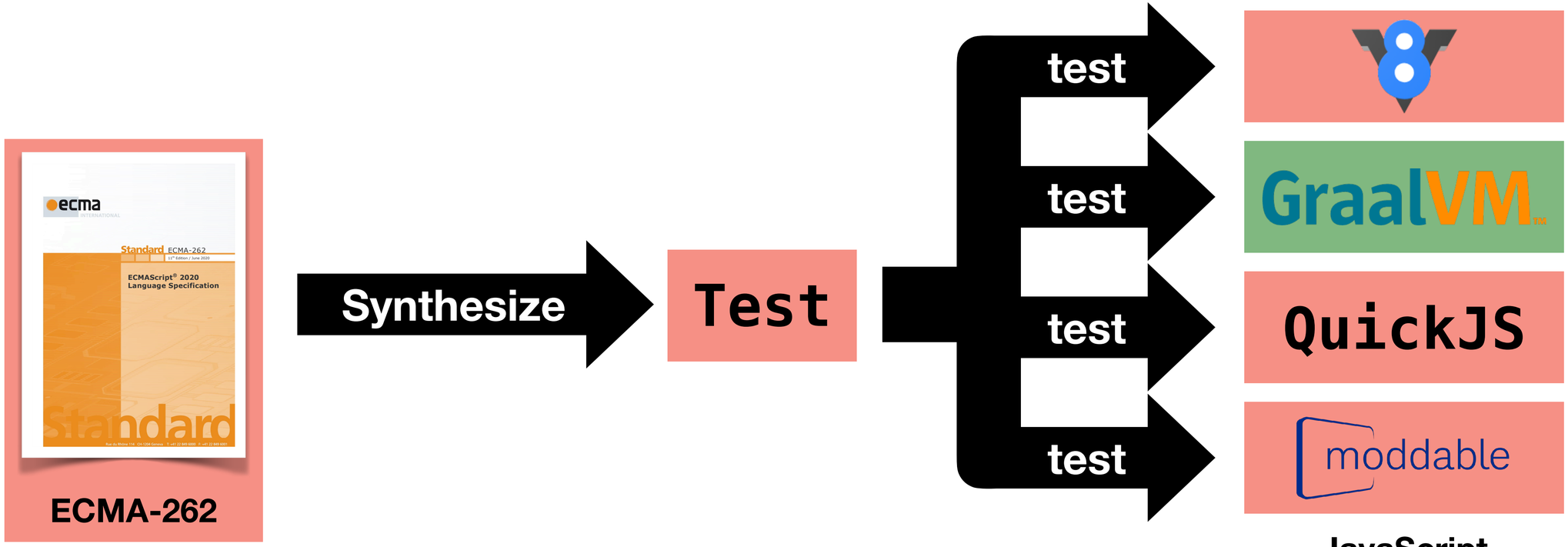


JEST - N+1-version Differential Testing



A <u>specification</u> bug in ECMA-262

JEST - N+1-version Differential Testing

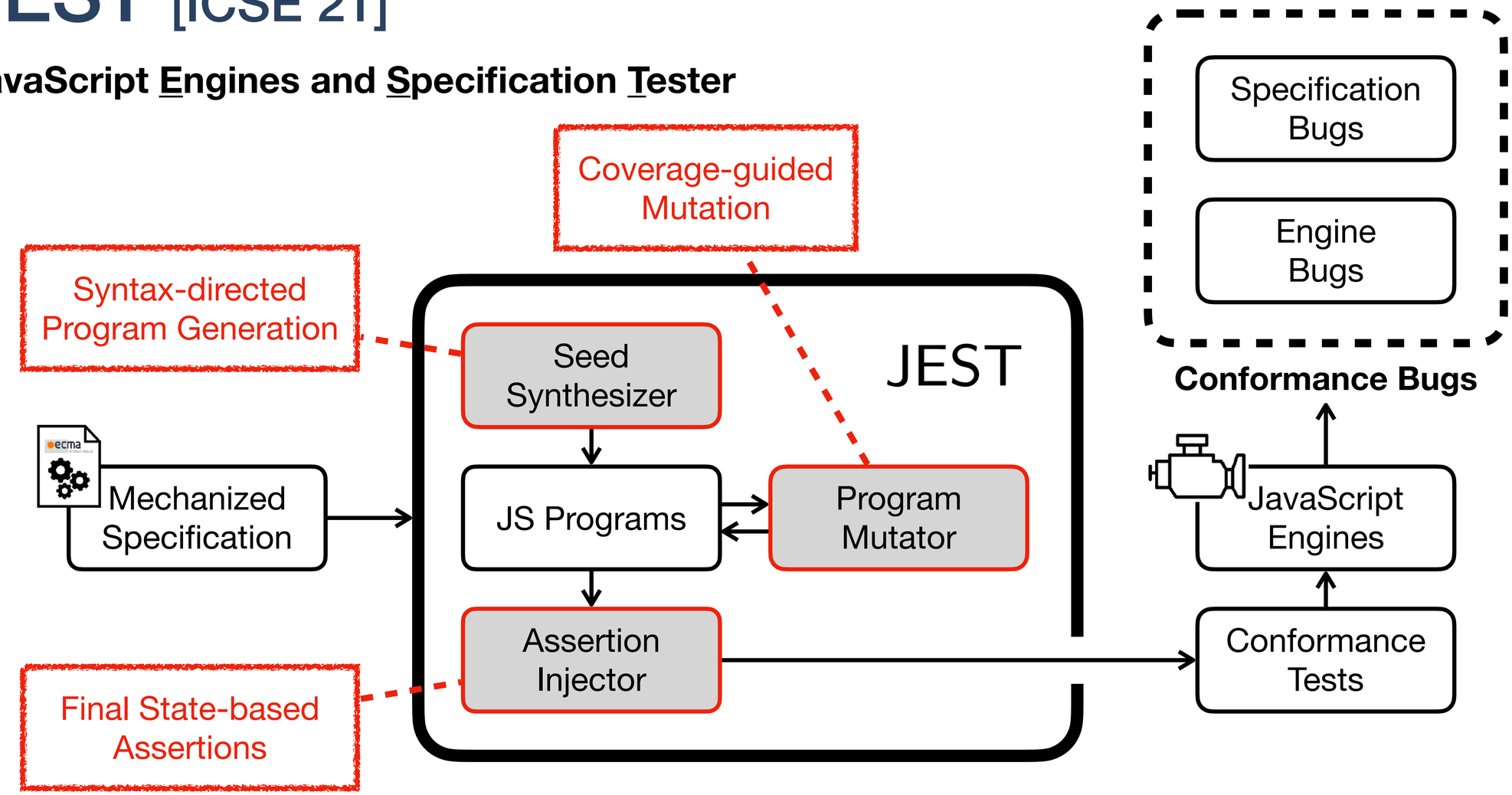


A <u>specification</u> bug in ECMA-262

An <u>engine</u> bug in GraalVM

JEST [ICSE'21]

JavaScript Engines and Specification Tester



JEST - Assertion Injector (7 Kinds)

```
var x = 1 + 2;
```

JEST - Assertion Injector (7 Kinds)

```
var x = 1 + 2;

+ $assert.sameValue(x, 3);
```

JEST - Assertion Injector (7 Kinds)

1. **Exceptions** (`Exc`)

```
+ // Throw
  let x = 42;
  function x() {};
```

2. **Aborts** (`Abort`)

```
+ // Abort
  var x = 42; x++;
```

3. **Variable Values** (`Var`)

```
  var x = 1 + 2;
+ $assert.sameValue(x, 3);
```

4. **Object Values** (`Obj`)

```
  var x = {}, y = {}, z = { p: x, q: y };
+ $assert.sameValue(z.p, x);
+ $assert.sameValue(z.q, y);
```

JEST - Assertion Injector (7 Kinds)

5. **Object Properties (`Desc`)**

```
  var x = { p: 42 };
+ $verifyProperty(x, "p", {
+   value: 42.0, writable: true,
+   enumerable: true, configurable: true
+ });
```

6. **Property Keys (`Key`)**

```
  var x = {[Symbol.match]: 0, p: 0, 3: 0, q: 0, 1: 0}
+ $assert.compareArray(
+   Reflect.ownKeys(x),
+   ["1", "3", "p", "q", Symbol.match]
+ );
```

7. **Internal Methods and Slots (`In`)**

```
  function f() {}
+ $assert.sameValue(Object.getPrototypeOf(f),
+                   Function.prototype);
+ $assert.sameValue(Object.isExtensible(x), true);
+ $assert.callable(f);
+ $assert.constructable(f);
```

JEST - Evaluation

44 Bugs in Engines

TABLE II: The number of engine bugs detected by JEST

Engines	Exc	Abort	Var	Obj	Desc	Key	In	Total
V8	0	0	0	0	0	2	0	2
GraalJS	6	0	0	0	2	8	0	16
QuickJS	3	0	1	0	0	2	0	6
Moddable XS	12	0	0	0	3	5	0	20
Total	21	0	1	0	5	17	0	44

27 Bugs in Spec.

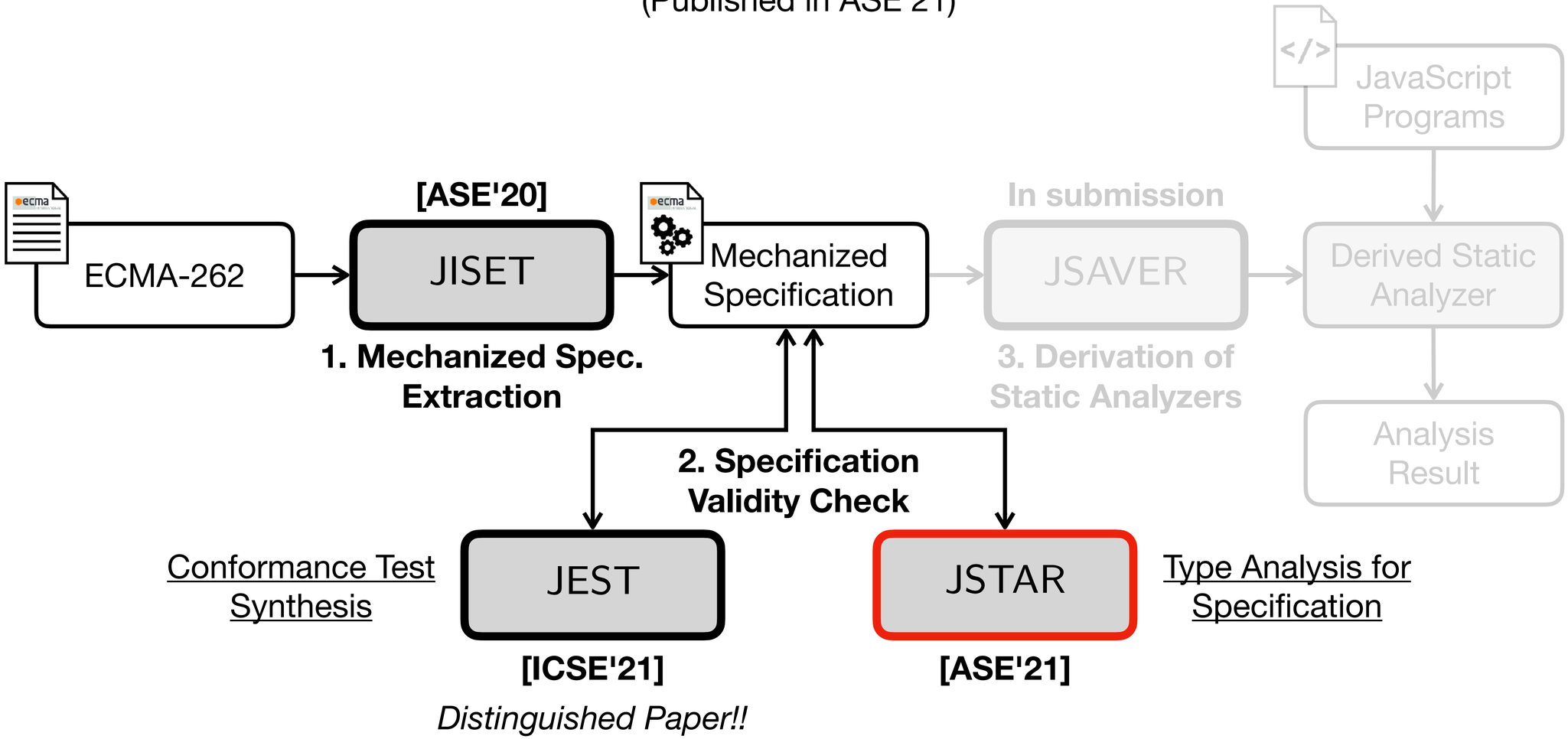
TABLE III: Specification bugs in ECMAScript 2020 (ES11) detected by JEST

Name	Feature	#	Assertion	Known	Created	Resolved	Existed
ES11-1	Function	12	Key	O	2019-02-07	2020-04-11	429 days
ES11-2	Function	8	Key	O	2015-06-01	2020-04-11	1,776 days
ES11-3	Loop	1	Exc	O	2017-10-17	2020-04-30	926 days
ES11-4	Expression	4	Abort	O	2019-09-27	2020-04-23	209 days
ES11-5	Expression	1	Exc	O	2015-06-01	2020-04-28	1,793 days
ES11-6	Object	1	Exc	X	2019-02-07	2020-11-05	637 days

JSTAR: JavaScript Specification Type Analyzer using Refinement

Jihyeok Park, Seungmin An, Wonho Shin, Yusung Sim, and Sukyoung Ryu

(Published in ASE'21)



JSTAR - Types in Specification

20.3.2.28 Math.round (x)

1. Let n be ? ToNumber(x).
2. If n is an integral Number, return n.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq$ -0.5, return **-0**.

...

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR - Types in Specification

20.3.2.28 Math.round (x **)** `x: (String v Boolean v Number v Object v ...)`

1. Let n be ? ToNumber(x).
2. If n is an integral Number, return n.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq -0.5$, return **-0**.

...

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR - Types in Specification

20.3.2.28 Math.round (x **)** `x: (String v Boolean v Number v Object v ...)`

1. Let n be ? ToNumber(x). `ToNumber(x): (Number v Exception)`
2. If n is an integral Number, return n.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq -0.5$, return **-0**.

...

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR - Types in Specification

20.3.2.28 Math.round (x) `x: (String v Boolean v Number v Object v ...)`

1. Let n be ? ToNumber(x). `ToNumber(x): (Number v Exception) ∧ n: (Number)`
2. If n is an integral Number, return n.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq -0.5$, return **-0**.

...

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR - Types in Specification

20.3.2.28 Math.round (x) `x: (String v Boolean v Number v Object v ...)`

1. Let n be ? ToNumber(x). `ToNumber(x): (Number v Exception) ∧ n: (Number)`
2. If n is an integral Number, return n.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq -0.5$, return **-0**.

...

Type Mismatch for numeric operator `>`

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR - Types in Specification

20.3.2.28 Math.round (x) `x: (String v Boolean v Number v Object v ...)`

1. Let n be ? ToNumber(x). `ToNumber(x): (Number v Exception) ∧ n: (Number)`
2. If n is an integral Number, return n.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq -0.5$, return **-0**.

...

Type Mismatch for numeric operator `>`

```
Math.round(true)  = ???
Math.round(false) = ???
```

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR - Types in Specification

20.3.2.28 Math.round (*x*) `x: (String v Boolean v Number v Object v ...)`

1. Let *n* be ? ToNumber(*x*). `ToNumber(x): (Number v Exception) ∧ n: (Number)`
2. If *n* is an integral Number, return *n*.
3. If $x < 0.5$ and $x > 0$, return **+0**.
4. If $x < 0$ and $x \geq -0.5$, return **-0**.

...

Type Mismatch for numeric operator `>`

```
Math.round(true)  = ???
Math.round(false) = ???
```

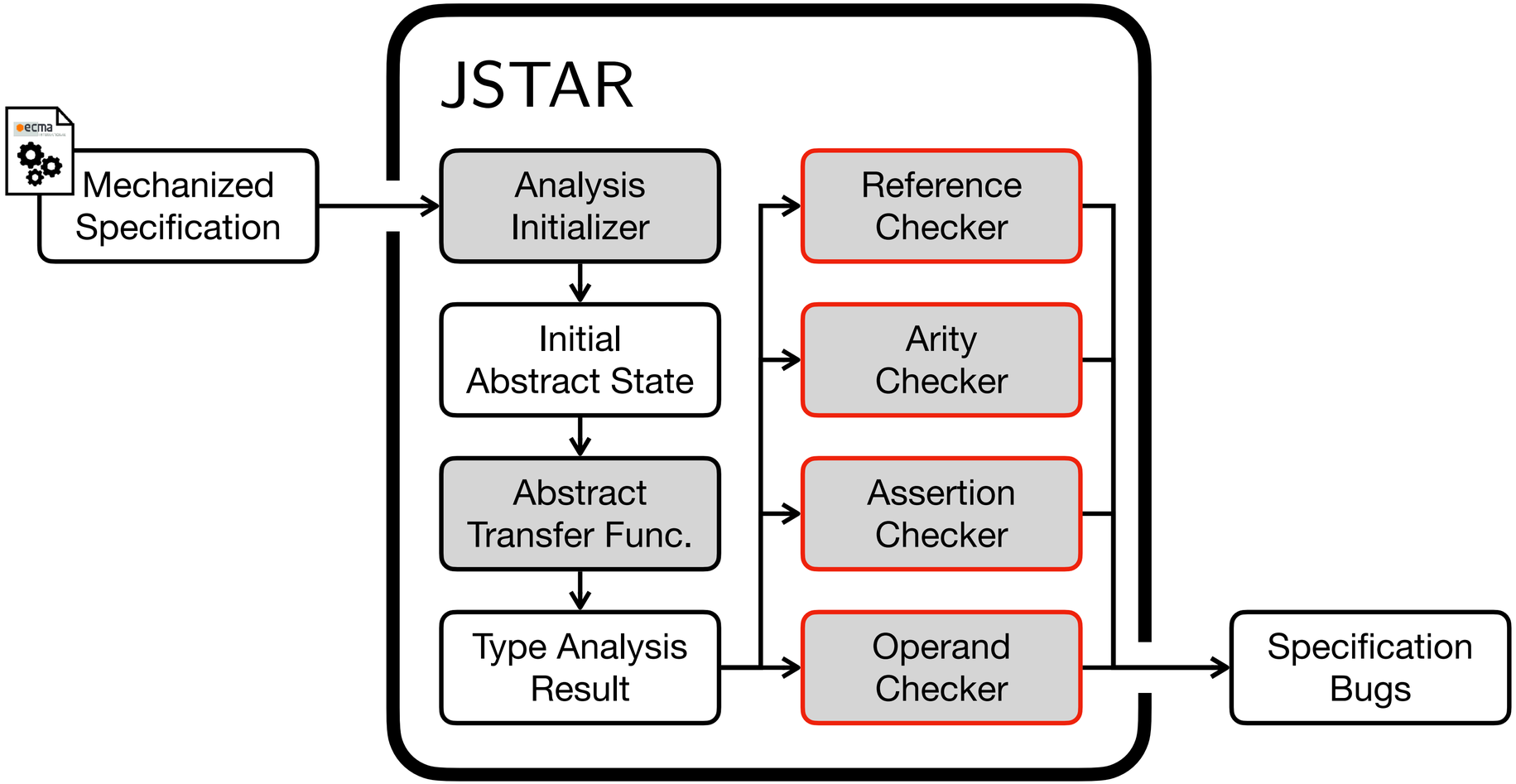
3. If $n < 0.5$ and $n > 0$, return **+0**.
4. If $n < 0$ and $n \geq -0.5$, return **-0**.

```
Math.round(true)  = 1
Math.round(false) = 0
```

https://github.com/tc39/ecma262/tree/575149cfd77aebcf3a129e165bd89e14caafc31c

JSTAR [ASE'21]

JavaScript Specification Type Analyzer using Refinement



JSTAR - Evaluation

- **Type Analysis for 864 versions of ECMA-262**

59.2% Precision

93 Bugs Detected

Checker	Bug Kind	Precision = (# True Bugs) / (# Detected Bugs)					
		no-refine		refine		Δ	
Reference	UnknownVar	62 / 106	17 / 60	63 / 78	17 / 31	+1 / -28	/ -29
	DuplicatedVar		45 / 46		46 / 47		+1 / +1
Arity	MissingParam	4 / 4	4 / 4	4 / 4	4 / 4	/	/
Assertion	Assertion	4 / 56	4 / 56	4 / 31	4 / 31	/ -25	/ -25
Operand	NoNumber	22 / 113	2 / 65	22 / 44	2 / 6	/ -69	/ -59
	Abrupt		20 / 48		20 / 38		/ -10
Total		92 / 279 (33.0%)		93 / 157 (59.2%)		+1 / -122 (+26.3%)	

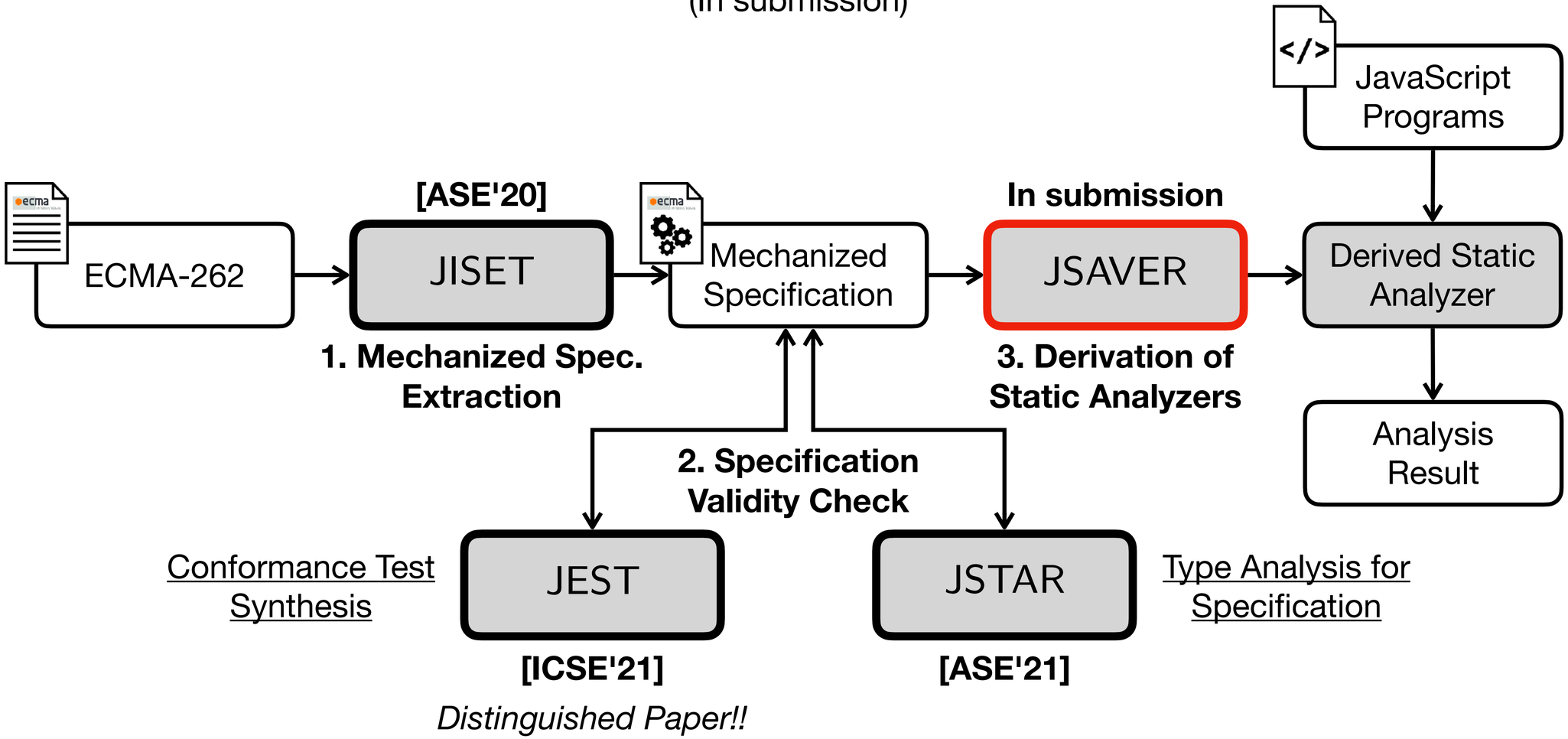
Name	Feature	#	Checker	Created	Life Span
ES12-1	Switch	3	Reference	2015-09-22	1,996 days
ES12-2	Try	3	Reference	2015-09-22	1,996 days
ES12-3	Arguments	1	Reference	2015-09-22	1,996 days
ES12-4	Array	2	Reference	2015-09-22	1,996 days
ES12-5	Async	1	Reference	2015-09-22	1,996 days
ES12-6	Class	1	Reference	2015-09-22	1,996 days
ES12-7	Branch	1	Reference	2015-09-22	1,996 days
ES12-8	Arguments	2	Operand	2015-12-16	1,910 days

14 Bugs in ES12

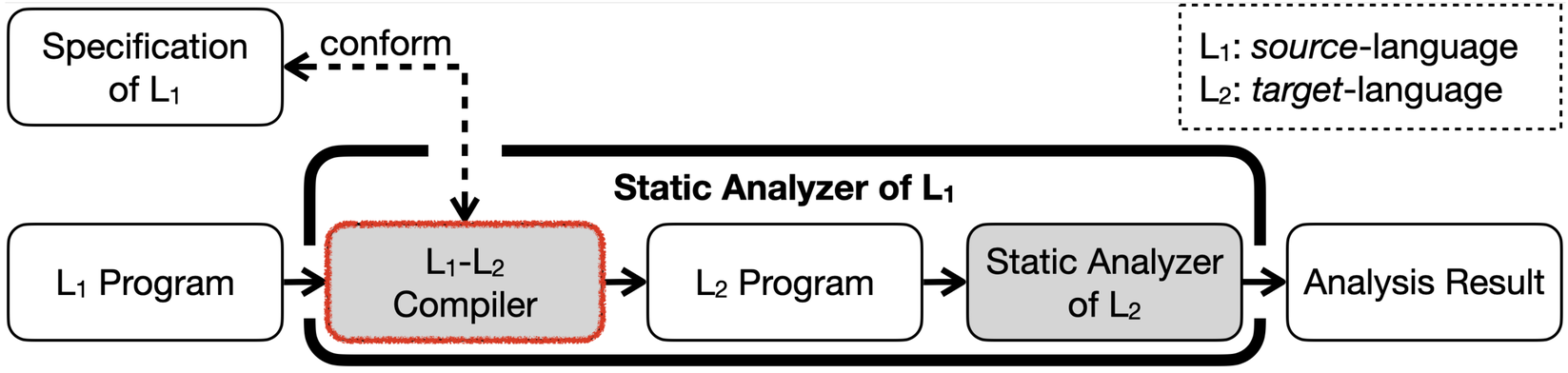
Automatically Deriving JavaScript Static Analyzers from Language Specifications

Jihyeok Park, Seungmin An, and Sukyoung Ryu

(In submission)

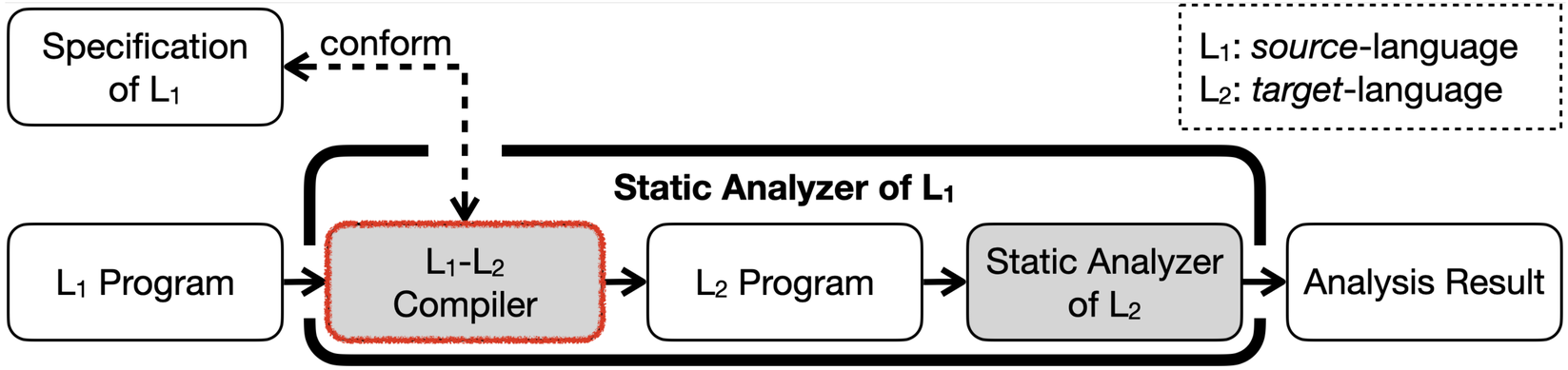


JSAVER - Meta-Level Static Analysis

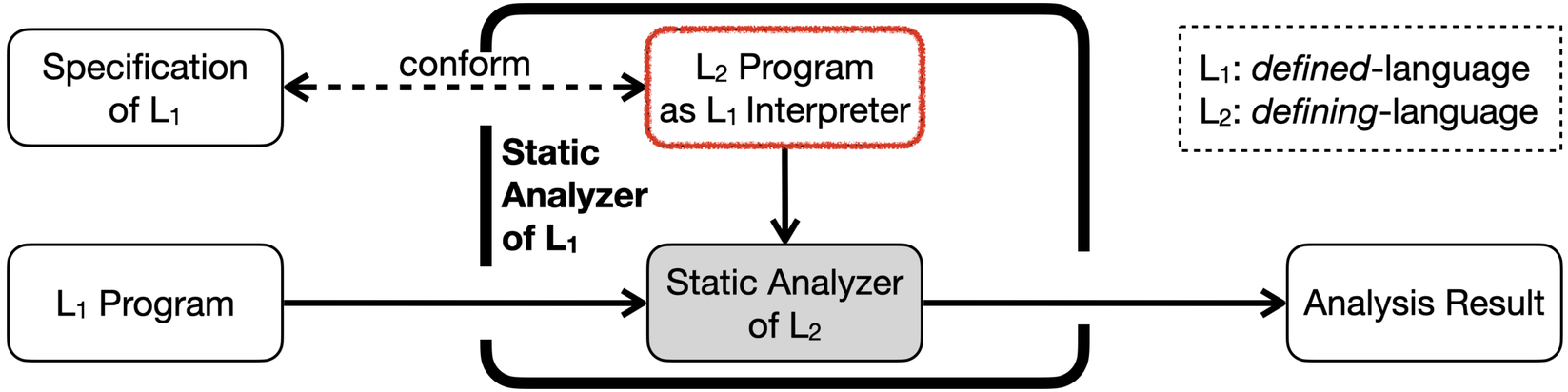


***compiler*-based approach (existing)**

JSAVER - Meta-Level Static Analysis



***compiler*-based approach (existing)**



***interpreter*-based approach (ours)**

JSAVER - Meta-Level Static Analysis

***defined*-language**
(JavaScript)

```
x ||= y
```

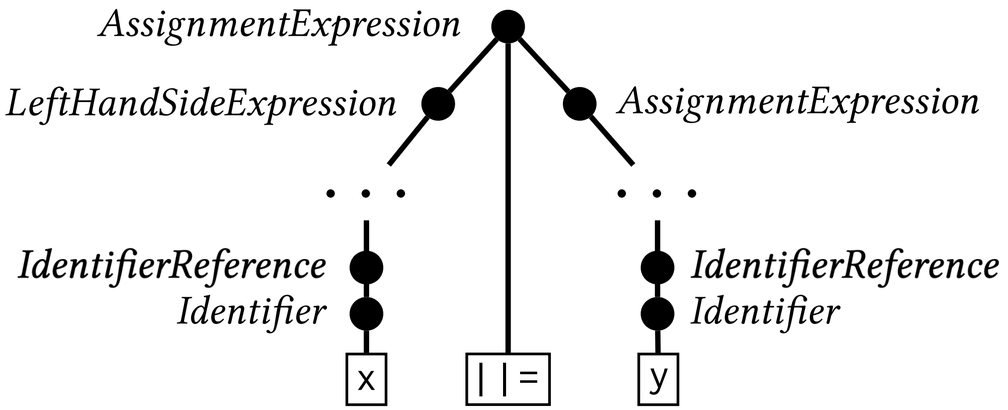
***defining*-language**
(IR_{ES})

JSAVER - Meta-Level Static Analysis

***defined*-language**
(JavaScript)

x ||= y

parse



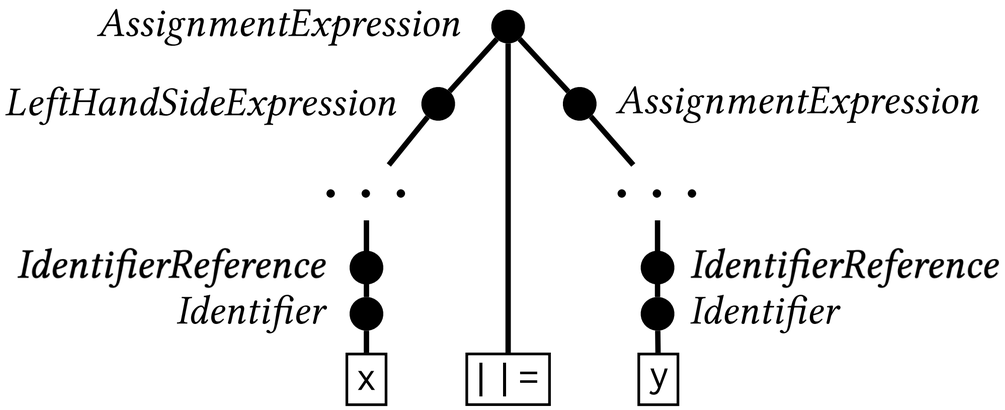
***defining*-language**
(IR$_{ES}$)

JSAVER - Meta-Level Static Analysis

***defined*-language**
(JavaScript)

x ||= y

parse



***defining*-language**
(IR_{ES})

```
syntax def AssignmentExpression[8].Evaluation(
  this, LeftHandSideExpression, AssignmentExpression
) {
  let lref = (LeftHandSideExpression.Evaluation)
  let lval = [? (GetValue lref)]
  let lbool = [! (ToBoolean lval)]
  if (= lbool true) return lval
  ...
}
```

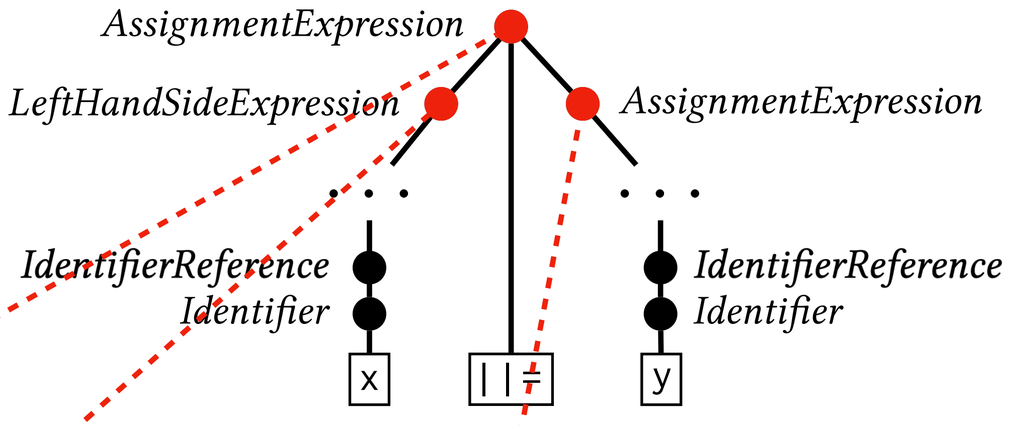
A mechanized specification from ES12
= A JavaScript interpreter
= An IR_{ES} program

JSAVER - Meta-Level Static Analysis

***defined*-language**
(JavaScript)

x ||= y

parse



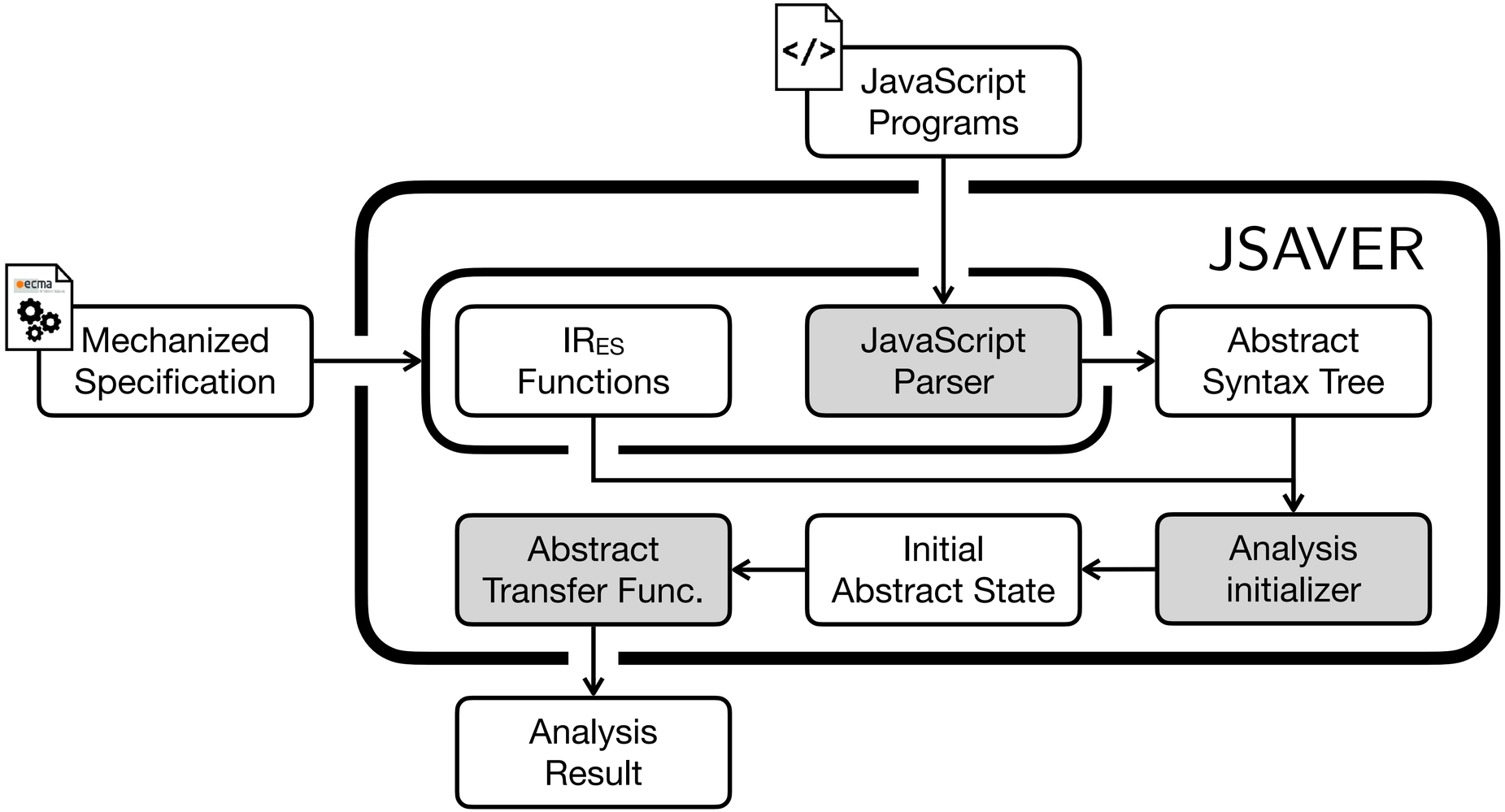
***defining*-language**
(IR_{ES})

```
syntax def AssignmentExpression[8].Evaluation(
  this, LeftHandSideExpression, AssignmentExpression
) {
  let lref = (LeftHandSideExpression.Evaluation)
  let lval = [? (GetValue lref)]
  let lbool = [! (ToBoolean lval)]
  if (= lbool true) return lval
  ...
}
```

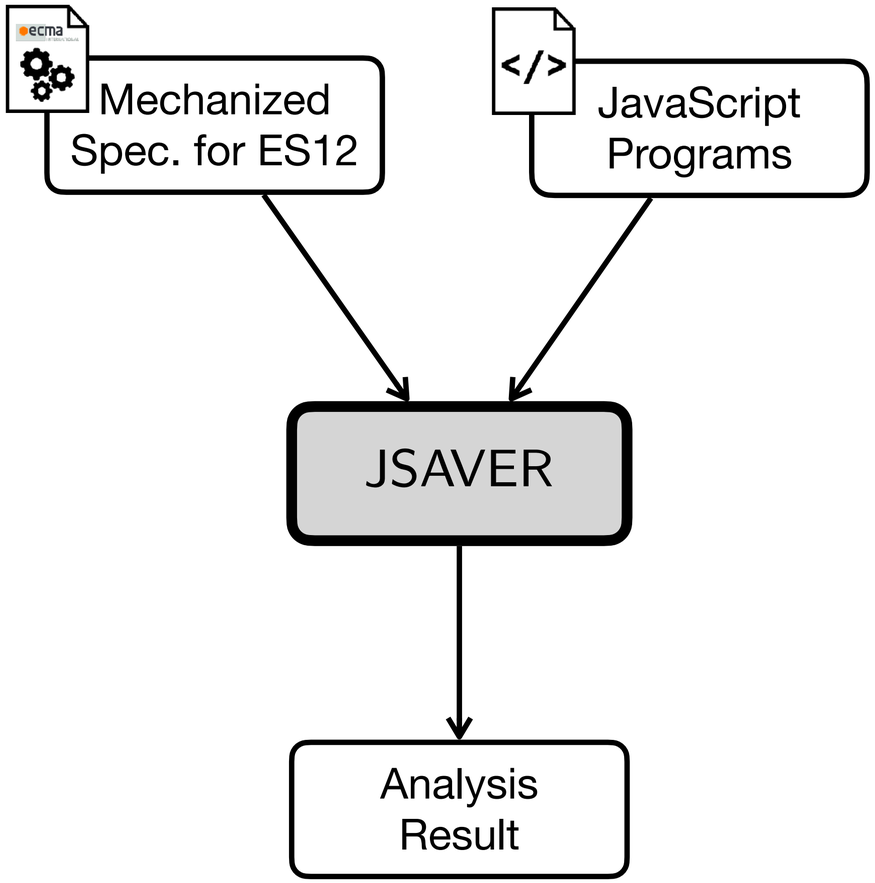
A mechanized specification from ES12
= A JavaScript interpreter
= An IR_{ES} program

JSAVER In submission

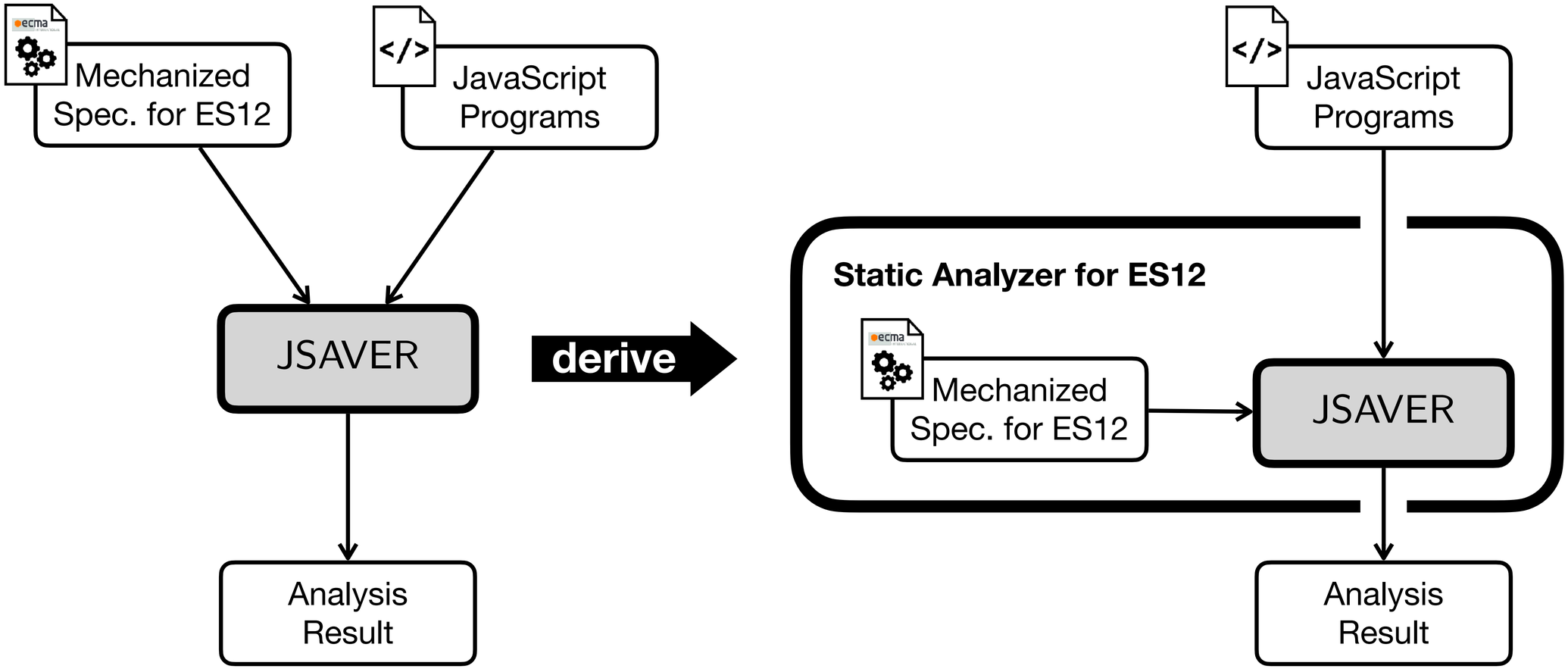
JavaScript Static Analyzer via ECMAScript Representation



JSAVER - Static Analyzer Derivation

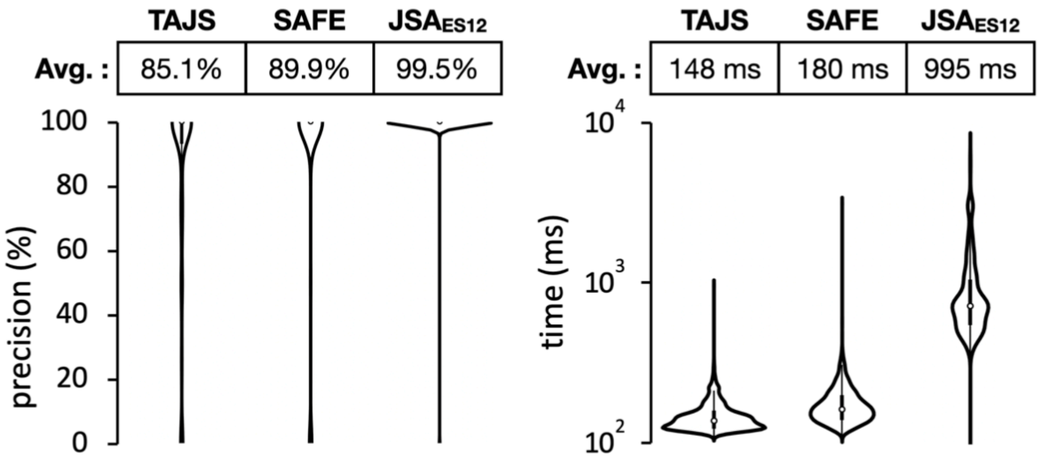


JSAVER - Static Analyzer Derivation

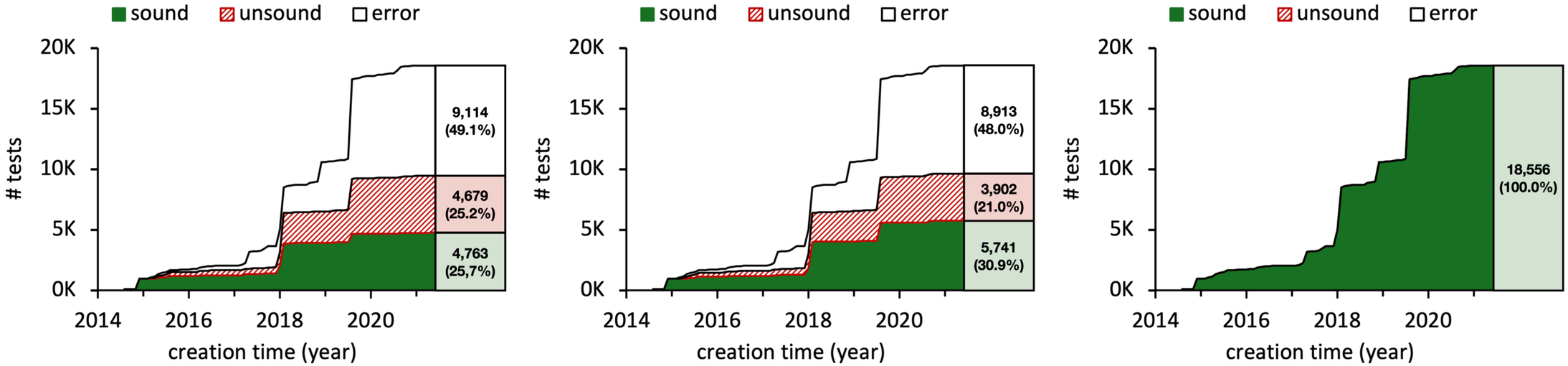


JSAVER - Evaluation

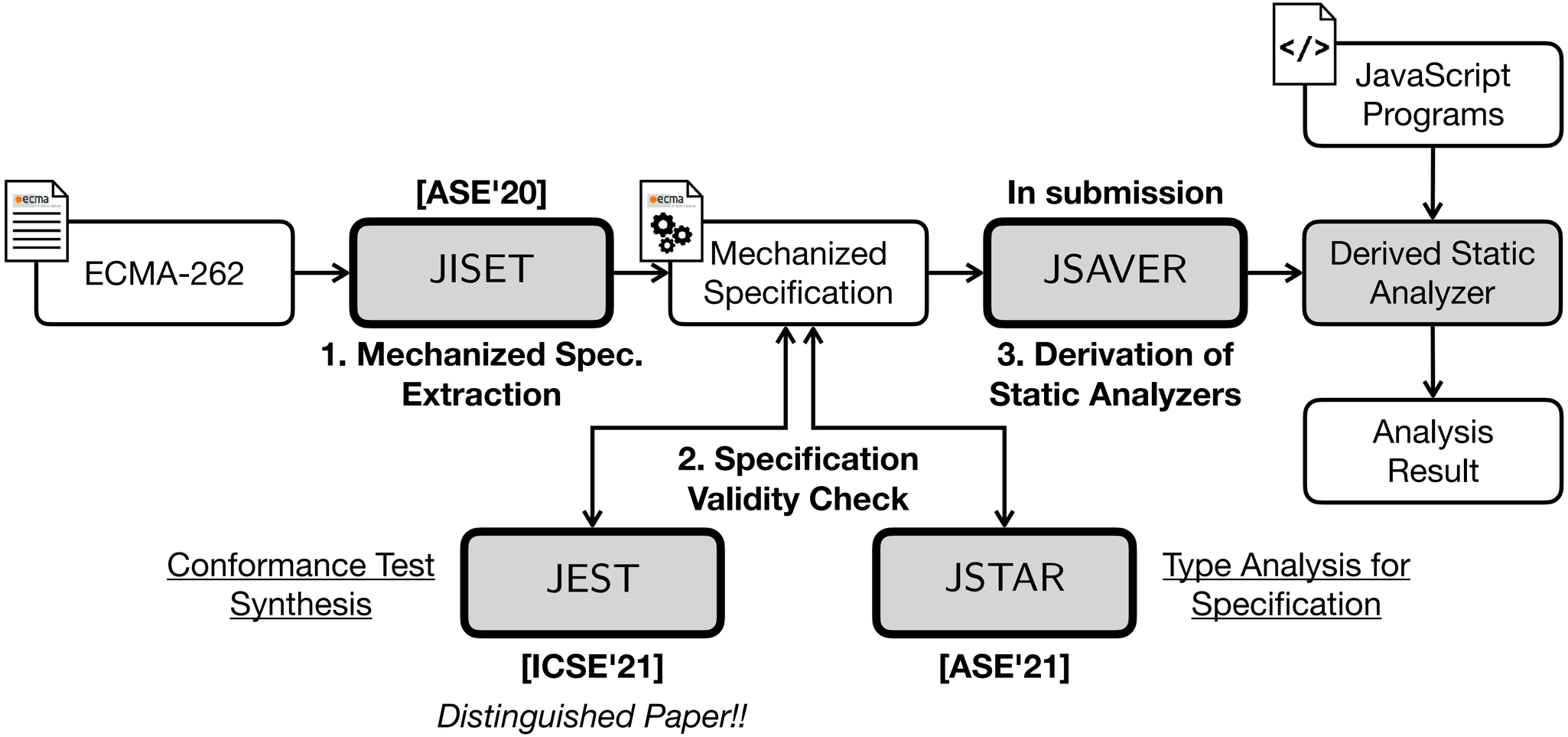
- **Soundness / Precision / Performance**
  - 18,556 applicable tests in Test262
  - 3,903 tests analyzable by all the three analyzers



(a) The analysis precision **(b)** The analysis performance



(a) Analysis results of TAJS **(b)** Analysis results of SAFE **(c)** Analysis results of JSA_{ES12}


JavaScript Programs
ECMA-262
[ASE'20]
JISET
Mechanized Specification
In submission
JSAVER
Derived Static Analyzer
1. Mechanized Spec. Extraction
3. Derivation of Static Analyzers
Analysis Result
2. Specification Validity Check
Conformance Test Synthesis
JEST
JSTAR
Type Analysis for Specification
[ICSE'21]
[ASE'21]
Distinguished Paper!!

ESMeta - es-meta/esmeta: ECMAScript Metalanguage for Generation of Language-based Tools

Search or jump to... / Pulls Issues Marketplace Explore

es-meta / esmeta Public

Pin | Unwatch 3 | Fork 0 | Starred 1

Code | Issues 11 | Pull requests | Discussions | Actions | Projects | ...

main | Go to file | Add file | Code

jhnaldo Removed unnecessary leg... ... 38 minutes ago 233

.github/workfl...	Update ci.yml	15 days ago
project	More project setting	20 days ago
src	Removed unnecessary legacy files	38 minutes ago
tests	Removed unnecessary legacy files	38 minutes ago
.gitignore	Added tests for grammar Stringifi...	15 days ago
.jvmopts	Added .jvmopts	14 days ago
.scalafmt.conf	Reformatted code with modified s...	2 days ago
LICENSE	sbt project setting	21 days ago

About

ECMAScript Metalanguage for Generation of Language-based Tools

Readme

BSD-3-Clause License

1 star

3 watching

0 forks

Releases

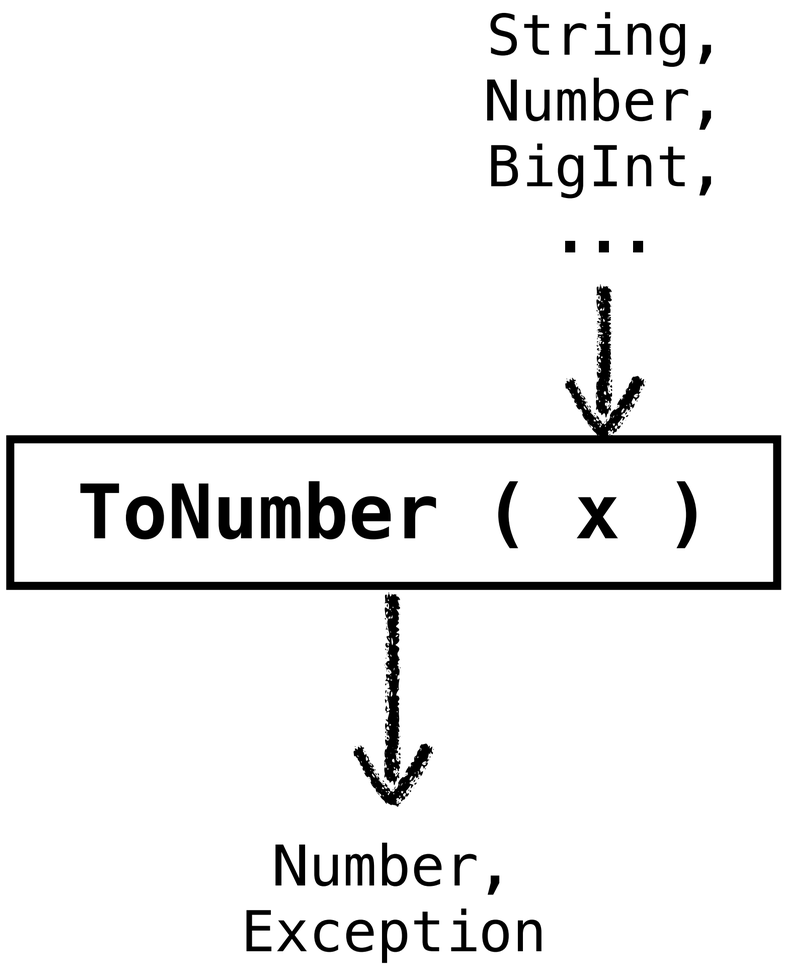
No releases published

Create a new release

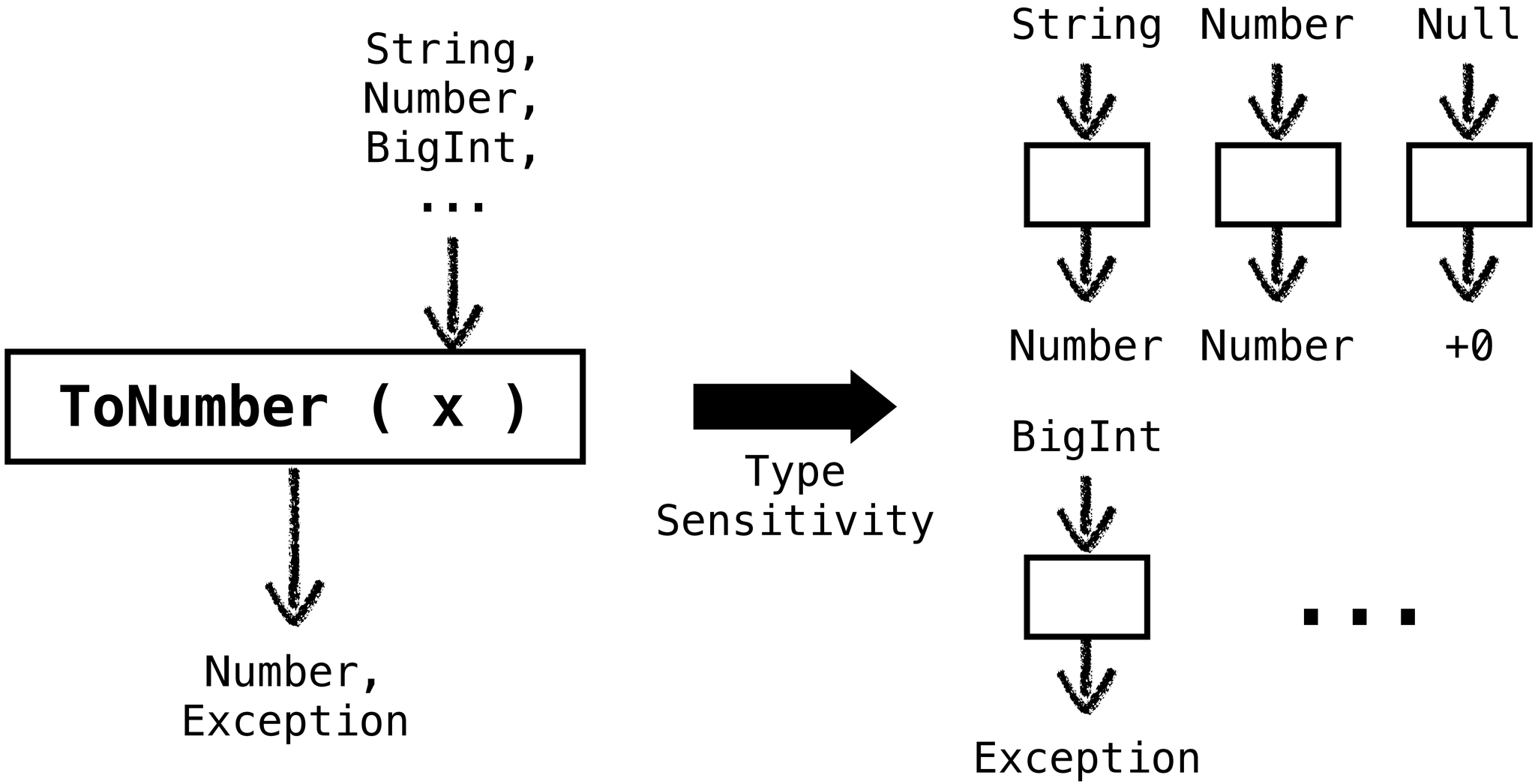
https://github.com/es-meta/esmeta

Backup Slides

JSTAR - Precision ⇧ - 1) Type Sensitivity



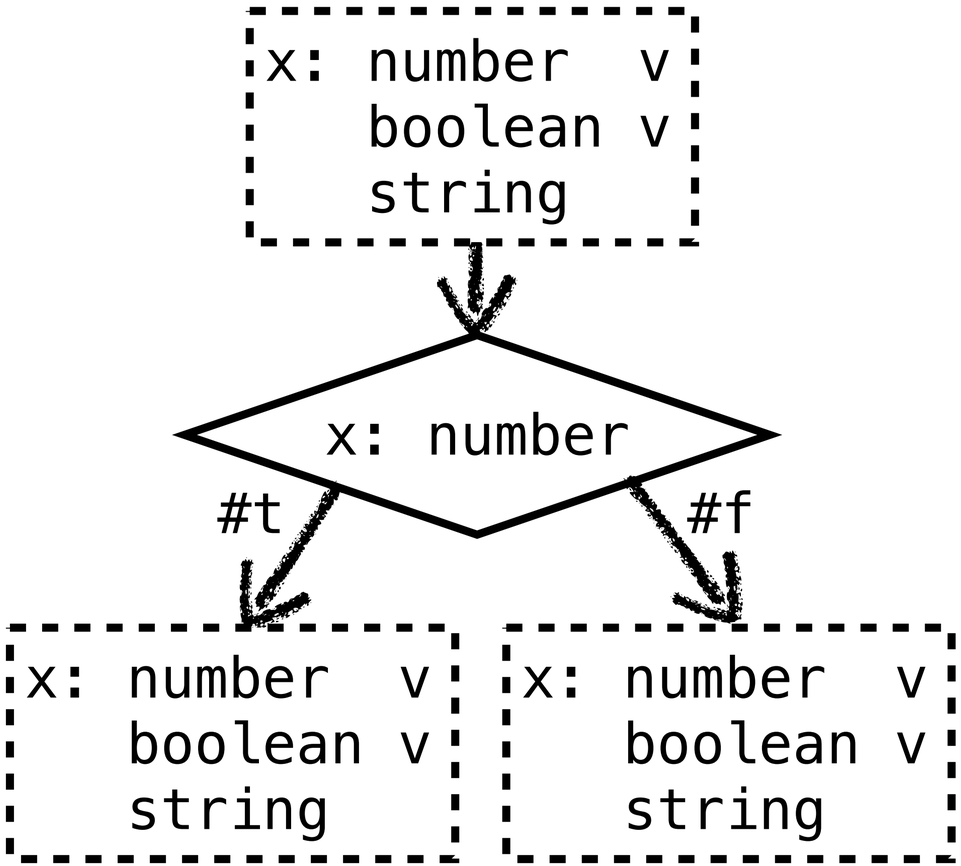
JSTAR - Precision ⇧ - 1) Type Sensitivity



JSTAR - Precision ⇧ - 2) Type Refinement

$$
\begin{aligned}
\texttt{refine}(!e, b)(\sigma^\sharp) &= \texttt{refine}(e, \neg b)(\sigma^\sharp) \\
\texttt{refine}(e_0 \texttt{ || } e_1, b)(\sigma^\sharp) &= \begin{cases} \sigma_0^\sharp \sqcup \sigma_1^\sharp & \text{if } b \\ \sigma_0^\sharp \sqcap \sigma_1^\sharp & \text{if } \neg b \end{cases} \\
\texttt{refine}(e_0 \texttt{ \&\& } e_1, b)(\sigma^\sharp) &= \begin{cases} \sigma_0^\sharp \sqcap \sigma_1^\sharp & \text{if } b \\ \sigma_0^\sharp \sqcup \sigma_1^\sharp & \text{if } \neg b \end{cases} \\
\texttt{refine}(\texttt{x.Type} == c_{\texttt{normal}}, \texttt{\#t})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \texttt{normal}(\mathbb{T})] \\
\texttt{refine}(\texttt{x.Type} == c_{\texttt{normal}}, \texttt{\#f})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \{\texttt{abrupt}\}] \\
\texttt{refine}(\texttt{x} == e, \texttt{\#t})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \tau_e^\sharp] \\
\texttt{refine}(\texttt{x} == e, \texttt{\#f})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \setminus \lfloor \tau_e^\sharp \rfloor] \\
\texttt{refine}(\texttt{x} : \tau, \texttt{\#t})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \{\tau\}] \\
\texttt{refine}(\texttt{x} : \tau, \texttt{\#f})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \setminus \{\tau' \mid \tau' <: \tau\}] \\
\texttt{refine}(e, b)(\sigma^\sharp) &= \sigma^\sharp
\end{aligned}
$$

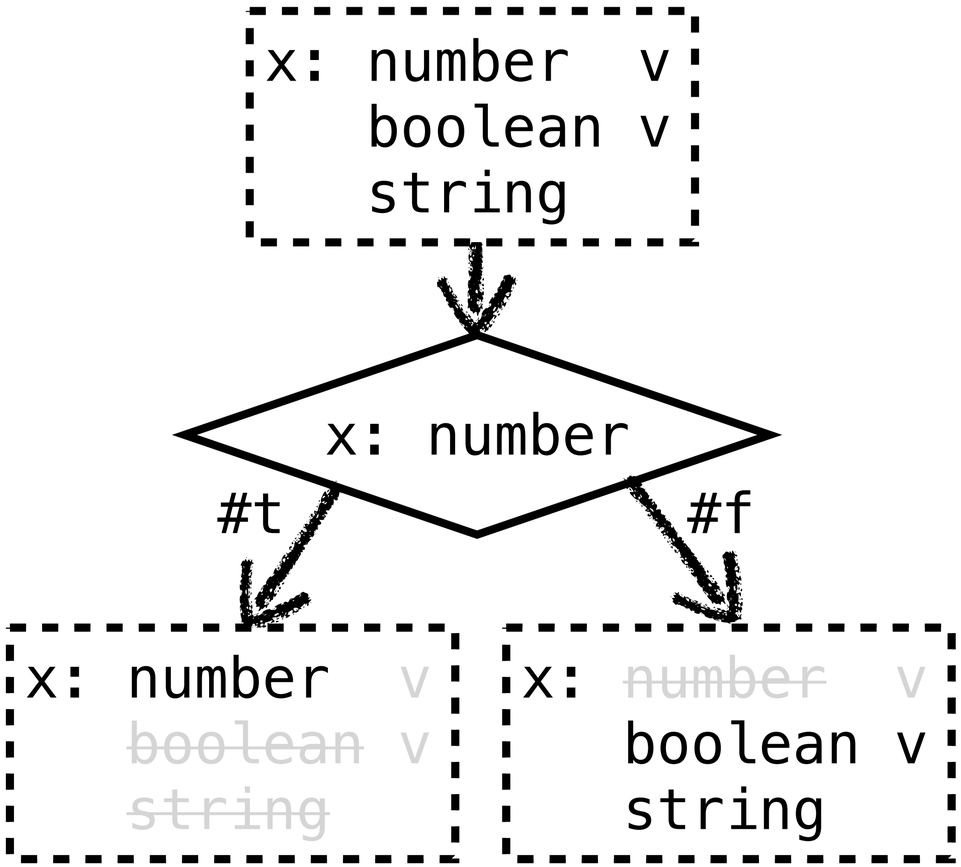
where $\sigma_j^\sharp = \texttt{refine}(e_j, b)(\sigma^\sharp)$ for $j = 0, 1$, $\tau_e^\sharp = [\![e]\!]_e^\sharp(\sigma^\sharp)$, and $\lfloor \tau^\sharp \rfloor$ returns $\{\tau\}$ if $\tau^\sharp$ denotes a singleton type τ, or returns $\varnothing$, otherwise.



JSTAR - Precision ⇧ - 2) Type Refinement

$$
\begin{aligned}
\texttt{refine}(!e, b)(\sigma^\sharp) &= \texttt{refine}(e, \neg b)(\sigma^\sharp) \\
\texttt{refine}(e_0 \texttt{ || } e_1, b)(\sigma^\sharp) &= \begin{cases} \sigma_0^\sharp \sqcup \sigma_1^\sharp & \text{if } b \\ \sigma_0^\sharp \sqcap \sigma_1^\sharp & \text{if } \neg b \end{cases} \\
\texttt{refine}(e_0 \texttt{ \&\& } e_1, b)(\sigma^\sharp) &= \begin{cases} \sigma_0^\sharp \sqcap \sigma_1^\sharp & \text{if } b \\ \sigma_0^\sharp \sqcup \sigma_1^\sharp & \text{if } \neg b \end{cases} \\
\texttt{refine}(\texttt{x.Type} == c_{\texttt{normal}}, \texttt{\#t})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \texttt{normal}(\mathbb{T})] \\
\texttt{refine}(\texttt{x.Type} == c_{\texttt{normal}}, \texttt{\#f})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \{\texttt{abrupt}\}] \\
\texttt{refine}(\texttt{x} == e, \texttt{\#t})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \tau_e^\sharp] \\
\texttt{refine}(\texttt{x} == e, \texttt{\#f})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \setminus \lfloor \tau_e^\sharp \rfloor] \\
\texttt{refine}(\texttt{x} : \tau, \texttt{\#t})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \sqcap \{\tau\}] \\
\texttt{refine}(\texttt{x} : \tau, \texttt{\#f})(\sigma^\sharp) &= \sigma^\sharp[\texttt{x} \mapsto \tau_{\texttt{x}}^\sharp \setminus \{\tau' \mid \tau' <: \tau\}] \\
\texttt{refine}(e, b)(\sigma^\sharp) &= \sigma^\sharp
\end{aligned}
$$

where $\sigma_j^\sharp = \texttt{refine}(e_j, b)(\sigma^\sharp)$ for $j = 0, 1$, $\tau_e^\sharp = [\![e]\!]_e^\sharp(\sigma^\sharp)$, and $\lfloor \tau^\sharp \rfloor$ returns $\{\tau\}$ if $\tau^\sharp$ denotes a singleton type τ, or returns $\varnothing$, otherwise.

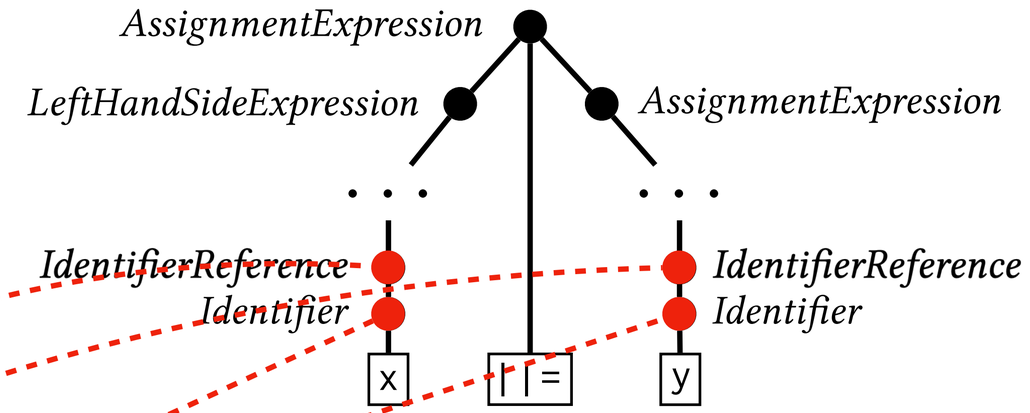


JSAVER - AST Sensitivity

***defined*-language**
(JavaScript)

x ||= y

parse



***defining*-language**
(IR$_{ES}$)

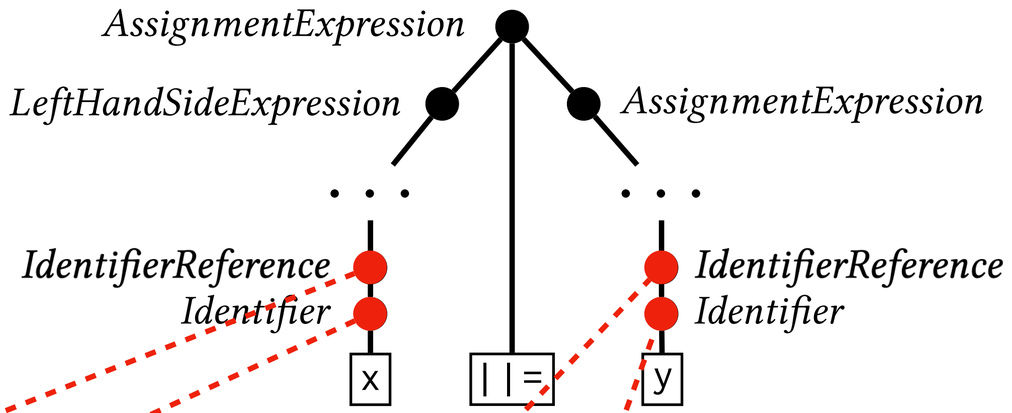
```
syntax def IdentifierReference[0]
  .Evaluation(
  this, Identifier
) {
  return [?
    (ResolveBinding
      (Identifier.StringValue))]
}
```

JSAVER - AST Sensitivity

***defined*-language**
(JavaScript)

```
x ||= y
```

parse



***defining*-language**
(IR$_{ES}$)

```
syntax def IdentifierReference[0]
  .Evaluation(
  this, Identifier
) {
  return [?
    (ResolveBinding
      (Identifier.StringValue))]
}
```

```
syntax def IdentifierReference[0]
  .Evaluation(
  this, Identifier
) {
  return [?
    (ResolveBinding
      (Identifier.StringValue))]
}
```

`this = AST of `x``

`this = AST of `y``

JSAVER - AST Sensitivity

defined-language (JavaScript)	**defining-language** ($\mathrm{IR}_{\mathrm{ES}}$)
flow-sensitivity	$\delta^{\text{js-flow}}(t_\perp) = \{\sigma = (\_, \_, \bar{c}, \_) \in \mathbb{S} \mid \mathrm{ast}(\bar{c}) = t_\perp\}$
k-callsite sensitivity	$\delta^{\text{js-}k\text{-cfa}}([t_1, \cdots, t_n]) = \{\sigma = (\_, \_, \bar{c}, \_) \in \mathbb{S} \mid n \leq k \wedge (n = k \vee \text{js-ctxt}^{n+1}(\bar{c}) = \perp) \wedge \forall 1 \leq i \leq n.\ \mathrm{ast} \circ \text{js-ctxt}^{i}(\bar{c}) = t_i\}$