Lecture 1 – Combinatorial Testing

AAA705: Software Testing and Quality Assurance

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PLRG

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Black-box Testing

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- Test data are derived from the **specification** of the software.
- In general, **exhaustive testing** is not feasible. It means that we **cannot guarantee** that the software is free of defects.
- We need to pick a good set of test cases to **maximize** the chance of **finding software errors**.

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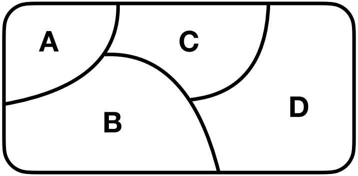
Equivalence Partitioning (EP)



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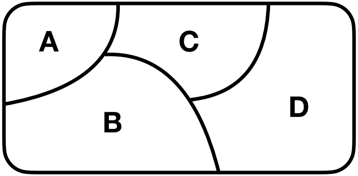

- **Equivalence partitioning** is a black-box testing technique that divides the input domain of a program into **equivalence classes**.
- The technique is based on the observation that the program should behave the same way for all members of an equivalence class.

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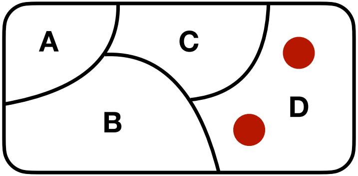

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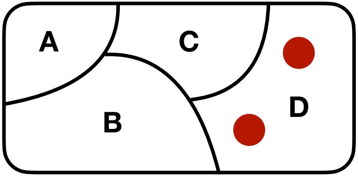

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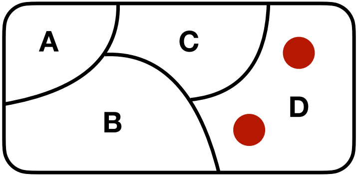

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- Then, how to define the equivalence classes?

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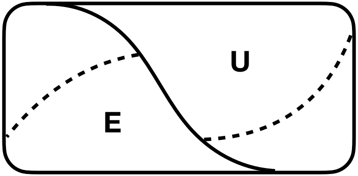

- One possible way to define the equivalence classes is to divide the input domain into **expected** and **unexpected** inputs.
  - **Expected (E)** or legal inputs
  - **Unexpected (U)** or illegal inputs

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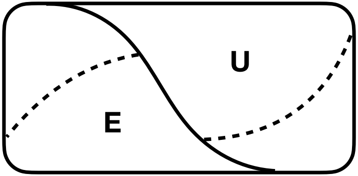

- One possible way to define the equivalence classes is to divide the input domain into **expected** and **unexpected** inputs.
  - **Expected (E)** or legal inputs
  - **Unexpected (U)** or illegal inputs

- We can further divide the expected inputs into smaller equivalence classes.

Equivalence Partitioning (EP)

Example

Consider a program that takes a **password** as input. The length of the password must be between 6 and 20 characters.

- We can divide the input domain into two equivalence classes:
  - $E = \{ \text{ a password } p \mid 6 \leq |p| \leq 20\}$
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- We can divide it more finely:
  - $E_1 = \{ \text{ a password } p \mid 6 \leq |p| \leq 10\}$ for *weak passwords*
  - $E_2 = \{ \text{ a password } p \mid 11 \leq |p| \leq 15\}$ for *medium-strength passwords*
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- We can select one test case from each equivalence class.

$$I = \{p_1, p_2, p_3, p_4, p_5\}$$

such that $|p_1| = 7$, $|p_2| = 13$, $|p_3| = 18$, $|p_4| = 3$, $|p_5| = 40$

Equivalence Partitioning (EP)

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- There are **many ways** to partition the input domain.

[1] [TSE'91] E. J. Weyuker and B. Jeng, "Analyzing partition testing strategies"

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- **Effectiveness** may depend on the tester's **experience** and **intuition**.

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Equivalence Partitioning (EP)

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  Partition testing can be better, worse, or the same as random testing, depending on how the partitioning is done.[1]

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- They usually occur due to **off-by-one** errors caused by misunderstanding the **boundary conditions**.
- It is **simple** but actually **very common**.

Off-by-one Errors – Looping Over Arrays

```
for (let i = 0; i < 10; i++) {
  /* body of the loop */
}
```

CORRECT

```
for (let i = 1; i < 10; i++) {
  /* body of the loop */
}
```

INCORRECT

```
for (let i = 0; i <= 10; i++) {
  /* body of the loop */
}
```

INCORRECT

```
for (let i = 0; i < 11; i++) {
  /* body of the loop */
}
```

INCORRECT

Off-by-one Errors – Fencepost error

If you build a straight fence 15 meters long with posts spaced 3 meters apart, how many posts do you need?

$15\ /\ 3 = 5$ posts?

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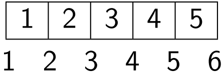

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```
linspace(0, 10, 5) == [0, 2.5, 5, 7.5, 10]
                   != [0, 2, 4, 6, 8, 10]
```

Off-by-one Errors – strncat in C

```
void foo (char *s)
{
  char buf[15];
  memset(buf, 0, sizeof(buf));
  // Final parameter should be: sizeof(buf)-1
  strncat(buf, s, sizeof(buf));
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- Off-by-one errors are common in using the C library because it is **not consistent** with respect to whether one needs to subtract 1 byte.

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- For example, we need to subtract 1 byte from the length of the buffer in `strncat` but not in `fgets` or `strncpy`.

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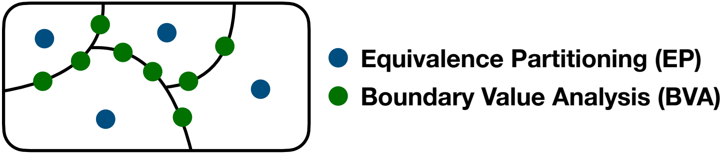
- For example, we need to subtract 1 byte from the length of the buffer in `strncat` but not in `fgets` or `strncpy`.
- So, the programmer has to remember for which functions they need to subtract 1.

Boundary Value Analysis (BVA)



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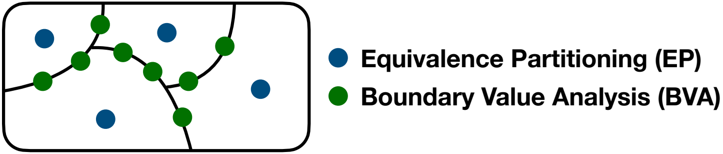

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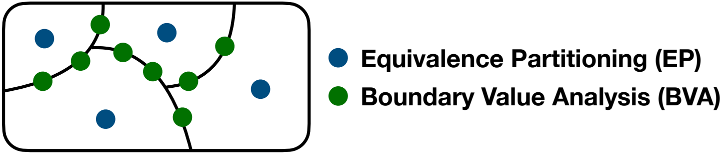

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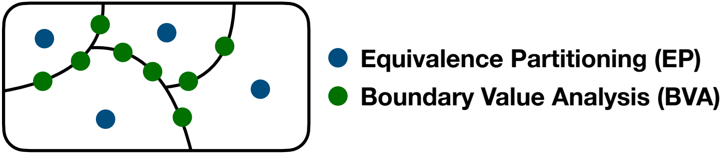

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- It is usually used in combination with **equivalence partitioning**.

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- We can select test cases at the boundaries of the equivalence classes.

$$I = \{p_1, p_2, p_3, p_4, p_5, p_6, p_7, p_8\}$$

such that

$$\begin{array}{llll} |p_1| = 5 & |p_2| = 6 & |p_3| = 10 & |p_4| = 11 \\ |p_5| = 15 & |p_6| = 16 & |p_7| = 20 & |p_8| = 21 \end{array}$$

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Category Partition Method (CPM)

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- Most programs behave differently when they receive different **parameters** or are executed under different **environments**.
- **Category partition method (CPM)** is a black-box testing technique that systematically generates test cases by considering the combinations of the categories of the input domain.
  1. Analyze **specification**
  2. Identify **parameters** and **environments**
  3. Identify **categories** for each parameter and environment
  4. **Partition** categories into equivalence classes
  5. Identify **constraints**
  6. **Generate** test cases

Category Partition Method (CPM)

Example

Unix command `grep` searches for files in a directory hierarchy with the following syntax:

```
grep <pattern> <filename>
```

For example,

- `grep park myfile` displays all lines in `myfile` that contain the word "`park`".
- `grep "hello world" myfile` displays all lines in `myfile` that contain the phrase "`hello world`".
- `grep " said \"hello " myfile` displays all lines in `myfile` that contain the phrase "` said "hello `".

Perform category partition method for the `grep` command.

Category Partition Method (CPM)

① Analyze **specification**

② Identify **parameters** and **environments**

- **Parameters** – (1) `<pattern>` and (2) `<filename>`
  - The `<pattern>` is a pattern to search for.
  - To include spaces in the pattern, it must be enclosed in quotes (`"`).
  - To include a quotation mark in the pattern, it must be escaped with a backslash (`\"`).
  - ...
- **Environments** – (3) file contents
  - ...

Category Partition Method (CPM)

③ Identify **categories** for each parameter and environment

1. **Parameter** – `<pattern>`
   - **Size**
   - **Quotation marks**
   - **Embedded spaces**
   - **Embedded quotation marks**
2. **Parameter** – `<filename>`
   - **Validity**
3. **Environment** – file contents
   - **Number of occurrences of the pattern**
   - **Number of occurrences of the pattern in a line**

Category Partition Method (CPM)

④ **Partition** categories into equivalence classes

1. **Parameter** – `<pattern>`
   - **Size** – 0 / 1 / ≥ 2
   - **Quotation marks** – quoted (Q) / unquoted (U) / improper (I)
   - **Embedded spaces** – none (N) / single (S) / multiple (M)
   - **Embedded quotation marks** – none (N) / single (S) / multiple (M)
2. **Parameter** – `<filename>`
   - **Validity** – exists (E) / not exists (N) / omitted (O)
3. **Environment** – file contents
   - **Number of occurrences of the pattern** – 0 / 1 / ≥ 2
   - **Number of occurrences of the pattern in a line** – 0 / 1 / ≥ 2

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How many combinations of the partitioned categories?

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⑤ Identify **constraints**

- For example, no embedded space for unquoted pattern.

$$\text{unquoted (U)} \quad \not\Rightarrow\!\Leftarrow \quad \text{single (S) for } \textbf{embedded spaces}$$
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⑥ **Generate** test cases

- Pick one **test case** from each combination **satisfying** the constraints.

Size	Q	Space	Emb. Q	Valid	Occur	Occur
0	Q	N	N	E	0	0
1	U	S	S	N	1	1
≥ 2	I	M	M	O	≥ 2	≥ 2

```
grep "hello world" myfile        # `myfile` is an empty file
```

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Airline	Destination	Departure Date	Return Date
79	171	365	365

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- **Combinatorial testing (CT)** or **combinatorial interaction testing (CIT)** constructs test cases by considering the **interactions** between the parameters.

Covering Array (CA)

Definition (Interaction)

For k parameters with v values each, a t-**way interaction** is a combination of values for t parameters.

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Definition (Covering Array (CA))

A **covering array** $A = CA(N; t, k, v)$ is a $N \times k$ matrix such that every field is an element from the set $[0, v-1]$, and every t-way interaction is covered at least once by a row of A.

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A	B	C
0	0	0
1	1	1

$\Rightarrow$

$CA(N = 4; t = 2, k = 3, v = 2)$

A	B	C
0	0	0
0	1	1
1	0	1
1	1	0

Mixed Covering Array (MCA)

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Definition (Mixed Covering Array (MCA))

A **mixed covering array** $A = CA(N; t, k, v = (v_1, v_2, \ldots, v_k))$ is a $N \times k$ matrix such that every field in the i-th column is an element from the set $[0, v_i - 1]$, and every t-way interaction is covered at least once by a row of A.

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A **mixed covering array** $A = CA(N; t, k, v = (v_1, v_2, \dots, v_k))$ is a $N \times k$ matrix such that every field in the i-th column is an element from the set $[0, v_i - 1]$, and every t-way interaction is covered at least once by a row of A.

A	B	C
0	0	0
1	1	1
		2

$\Rightarrow$

$CA(N = 6; t = 2, k = 3, v = (2, 2, 3))$

A	B	C
0	0	0
0	1	1
1	0	2
1	1	0
1	0	1
0	1	2

Constraint Mixed Covering Array (CMCA)

Definition (Constrinat Mixed Covering Array (CMCA))

A **mixed covering array** $A = CA(N; t, k, v = (v_1, v_2, \ldots, v_k), P)$ is a $N \times k$ matrix such that every field in the i-th column is an element from the set $[0, v_i - 1]$, and every **valid** t-way interaction is covered at least once by a row of A. We say that a t-way interaction is **valid** if it satisfies the predicate P.

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A	B	C
0	0	0
1	1	1
		2

$\Rightarrow$

$CA(N = 5; t = 2, k = 3, v = (2, 2, 3), P)$

A	B	C
0	1	1
1	0	2
1	1	0
1	0	1
0	1	2

$$P(x, y) = x + y > 0$$

Combinatorial Testing (CT)

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Pairwise testing is a special case of combinatorial testing with $t = 2$.

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Definition (Pairwise Testing)

Pairwise testing is a special case of combinatorial testing with $t = 2$.

Pairwise testing produces 5 test cases for the following system:

A	B	C
0	0	0
1	1	1
		2

$$P(x, y) = x + y > 0$$

$\Rightarrow$

$$CA(N = 5; t = 2, k = 3, v = (2, 2, 3), P)$$

A	B	C
0	1	1
1	0	2
1	1	0
1	0	1
0	1	2

Combinatorial Testing (CT)

Definition (Combinatorial Testing (CT))

Combinatorial testing with a strength t produces a test suite from a covering array $CA(N; t, k, v)$ for a system with k parameters, each with

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Definition (Pairwise Testing)

Pairwise testing is a special case of combinatorial testing with $t = 2$.

Pairwise testing produces 5 test cases for the following system:

A	B	C
0	0	0
1	1	1
		2

$$P(x, y) = x + y > 0$$

$\Rightarrow$

$$CA(N = 5; t = 2, k = 3, v = (2, 2, 3), P)$$

A	B	C
0	1	1
1	0	2
1	1	0
1	0	1
0	1	2

Fault Detection Effectiveness

- Pairwise testing is a **good trade-off** between test effort and test effectiveness.
  - For a system with 20 parameters each with 15 values, **pairwise testing** only requires 412 tests, whereas exhaustive testing requires $15^{20} = 3.5 \times 10^{25}$ tests.

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Fault Detection Effectiveness

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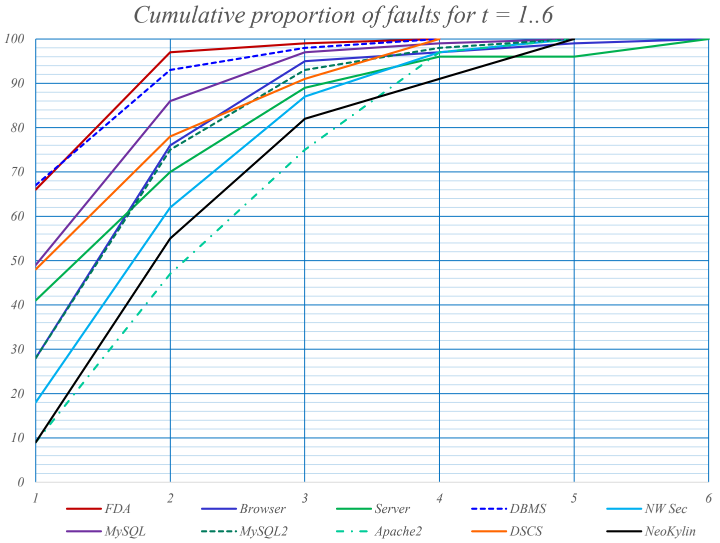
  - For a system with 20 parameters each with 15 values, **pairwise testing** only requires 412 tests, whereas exhaustive testing requires $15^{20} = 3.5 \times 10^{25}$ tests.
- Is higher strength always better for fault detection?
- It depends on the target program, but we can analyze the **general trend** against **a set of known faults**.
  - **Pairwise testing** discovers at least **53%** of the known faults.
  - **6-way testing** discovers **100%** of the known faults.

Fault Detection Effectiveness



"Combinatorial Methods in Software Testing" by Rick Kuhn, NIST,

Greedy Algorithm – IPOG Strategy

- The problem of generating a minimum covering array is **NP-complete**.
  - It can be reduced to the **vertex cover problem**.

[2] **[ECBS'07]** LEI, Yu, et al. "*IPOG: A general strategy for t-way software testing.*

Greedy Algorithm – IPOG Strategy

- The problem of generating a minimum covering array is **NP-complete**.
  - It can be reduced to the **vertex cover problem**.
- Let's learn a **polynomial time** greedy algorithm called **IPOG (In-Parameter-Order-General)**[2] that generates a covering array with a strength t. It is not optimal but practical.

[2] **[ECBS'07]** LEI, Yu, et al. "*IPOG: A general strategy for t-way software testing.*

Greedy Algorithm – IPOG Strategy

1. Initialize **test set** *ts* to be an empty set.
2. **Parameters** are P_1, P_2, ..., P_k.
3. Add a test into *ts* for **all interactions of the first t parameters**.
4. `for` ($i = t + 1$; $i \leq n$; i++) (**Horizontal Growth**)
   1. Let π be the **set of t-way interactions** involving **parameter** P_i and $t - 1$ parameters among the first $i - 1$ parameters.
   2. `for` (test $\gamma = (v_1, v_2, \cdots, v_{i-1}) \in ts$)
      1. **Choose** a value v_i of P_i
      2. **Replace** γ with $\gamma' = (v_1, v_2, \cdots, v_{i-1}, v_i)$ such that γ' covers the most number of interactions in π.
      3. **Remove** from π the interaction covered by γ'.
5. `for` (interaction $\alpha \in \pi$) (**Vertical Growth**)
   1. `if` ($\exists$ a test covers α) **Remove** α from π.
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Greedy Algorithm – IPOG Strategy

PLRG

Example

We want to **pairwise testing** for the following system:

P_1	P_2	P_3
0	0	0
1	1	1
		2

Adding all combinations of values between the first 2 parameters:

$ts =$

P_1	P_2
0	0
0	1
1	0
1	1

Greedy Algorithm – IPOG Strategy

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4. `for` $(i = t + 1;\ i \leq n;\ i{+}{+})$ (**Horizontal Growth**)
    1. Let π be the **set of t-way interactions** involving **parameter P_i** and $t - 1$ parameters among the first $i - 1$ parameters.

Greedy Algorithm – IPOG Strategy

Set π as pairs to cover involving P_3:

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

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        3. **Remove** from π the interaction covered by γ'.

Greedy Algorithm – IPOG Strategy

Adding values for P_3 in ts:

$ts =$

P_1	P_2	P_3
0	0	
0	1	
1	0	
1	1	

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

Adding values for P_3 in ts:

$ts =$

P_1	P_2	P_3
0	0	0
0	1	
1	0	
1	1	

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

Adding values for P_3 in ts:

$ts =$

P_1	P_2	P_3
0	0	0
0	1	1
1	0	
1	1	

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

PLRG

Adding values for P_3 in ts:

$ts =$

P_1	P_2	P_3
0	0	0
0	1	1
1	0	1
1	1	

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

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P_1	P_2	P_3
0		0
	0	0
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Greedy Algorithm – IPOG Strategy

Extending ts:

$ts =$

P_1	P_2	P_3
0	0	0
0	1	1
1	0	1
1	1	0

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

PLRG

Extending *ts*:

$ts =$

P_1	P_2	P_3
0	0	0
0	1	1
1	0	1
1	1	0
0	0	2

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy Algorithm – IPOG Strategy

PLRG

Extending *ts*:

$ts =$

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0	1	1
1	0	1
1	1	0
0	0	2
1	1	2

$\pi =$

P_1	P_2	P_3
0		0
	0	0
0		1
	0	1
0		2
	0	2
1		0
	1	0
1		1
	1	1
1		2
	1	2

Greedy vs. Meta-heuristic

- Simulated Annealing, a type of local search algorithm, has been proven to be effective against CIT.[3]

[3] **[SBSE'09]** B. J. Garvin et al. *"An improved meta-heuristic search for constrained interaction testing."*

Greedy vs. Meta-heuristic

- Simulated Annealing, a type of local search algorithm, has been proven to be effective against CIT.[3]
- The followings are size and time comparisons between the greedy algorithm and the meta-heuristic algorithm (average of 50 runs).

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Size comparison

Subject	Greedy	Meta-heuristic
SPIN-S	27	19
SPIN-V	42	36
GCC	24	21
Apache	42	32
Bugzilla	21	16

Time (sec.) comparison

Subject	Greedy	Meta-heuristic
SPIN-S	0.2	8.6
SPIN-V	11.3	102.1
GCC	204	1902.0
Apache	76.4	109.1
Bugzilla	1.9	9.1

[3] **[SBSE'09]** B. J. Garvin et al. *"An improved meta-heuristic search for constrained interaction testing."*

Summary

1. Equivalence Partitioning (EP)

2. Boundary Value Analysis (BVA)

3. Category Partition Method (CPM)

4. Combinatorial Testing (CT)
   - Covering Array (CA)
   - Fault Detection Effectiveness
   - Greedy Algorithm – IPOG Strategy
   - Greedy vs. Meta-heuristic

Next Lecture

- Random Testing

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