

Lecture 3 – Coverage Criteria

AAA705: Software Testing and Quality Assurance

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- Random Testing (RT)
 - Probabilistic Analysis
 - Weaknesses of Random Testing
 - Examples
- Adaptive Random Testing (ART)
 - Levenshtein (Edit) Distance
 - Distance Comparison Target
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1. Graph Coverage

- Structural Coverage

- Data-Flow Coverage

- Subsumption Relationships

2. Logic Coverage

- Simple Logic Expression Coverage

- Active Clause Coverage

- Inactive Clause Coverage

- Subsumption Relationships

3. Neuron Coverage

4. Feature-Sensitive Coverage

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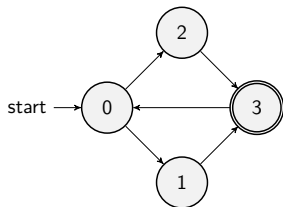
4. Feature-Sensitive Coverage

- **Graphs** are the most **commonly** used structure in software.
 - Control Flow Graphs (CFGs)
 - Call Graphs
 - Design Structure
 - Finite State Machines (FSMs)
 - etc.
- We want to ensure that our **tests** properly **cover** the graph.

Definition (Graph)

A **graph** $G = (N, E, N_s, N_f)$ is a quadruple consisting of

- 1 a set of **nodes** N ,
- 2 a set of **edges** $E \subseteq N \times N$,
- 3 a set of **start nodes** $N_s \subseteq N$, and
- 4 a set of **final nodes** $N_f \subseteq N$.



$$G = \begin{cases} N = \{0, 1, 2, 3\} \\ E = \{(0, 1), (0, 2), (1, 3), (2, 3), (3, 0)\} \\ N_s = \{0\} \\ N_f = \{3\} \end{cases}$$

- A **path** $p = (n_0, n_1, \dots, n_k) \in N^*$ in a graph $G = (N, E, N_s, N_F)$ is a sequence of nodes such that $(n_i, n_{i+1}) \in E$ for $0 \leq i < k$.

$$P_G = \{p \in N^* \mid p \text{ is a path in } G\}$$

- The **length** of a path is the **number of edges** in the path.

$$|p| = k$$

- A **subpath** q of a path p is a **subsequence** of p (i.e., $q \preceq p$).

$$q \preceq p \iff \exists 0 \leq i \leq j \leq k. q = (n_i, n_{i+1}, \dots, n_j)$$

- A path p is a **test path** if it starts from the **start node** and ends at a **final node**, and it represents an **execution** of a **test case**.

$$n_0 \in S \wedge n_k \in F$$

- A test path p **visits** a node n if it is in the path.

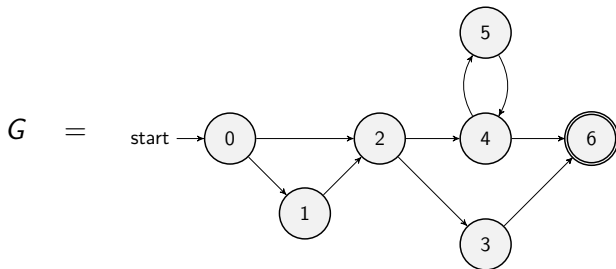
$$\exists 0 \leq i \leq k. n_i = n$$

- A test path p **visits** an edge (n, m) if it is in the path.

$$\exists 0 \leq i \leq k. n_i = n \wedge n_{i+1} = m$$

- The **path**(t) is the test path executed by a test case t .
- The **path**(T) is the set of test paths executed by a test suite T .

Consider a path $p = [0, 2, 4, 5, 4, 6]$ in the graph G .



- $|p| = 5$
- $[2, 4, 5]$ is a **subpath** of p (i.e., $[2, 4, 5] \preceq p$)
- p is a **test path**
- p **visits** nodes 0, 2, 4, 5, and 6
- p **visits** edges $(0, 2)$, $(2, 4)$, $(4, 5)$, $(5, 4)$, and $(4, 6)$

Definition (Graph Coverage Criterion)

A **graph coverage criterion** $C_G = (R_G, \sim_G)$ for a given graph G is defined with:

- a set of **test requirements (TRs)** R_G , and
- a **cover relation** $\sim \subseteq P_G \times R_G$ between paths and test requirements.
- A **test case** t **covers** a TR r if its test path satisfies the TR.

$$t \sim r \iff \text{path}(t) \sim r$$

- A **test suite** T **satisfies** the coverage criterion C_G if it covers all TRs.

$$T \vdash C_G \iff \forall r \in R_G. \exists t \in T. t \sim r$$

- A **structural coverage criterion** is defined on a graph in terms of **nodes**, **edges**, and **paths**.
- A **data-flow coverage criterion** is defined on a graph annotated with references to **variables**.

Definition (Node Coverage (NC))

The **node coverage** criterion $C_G = (R_G, \sim)$ is defined with:

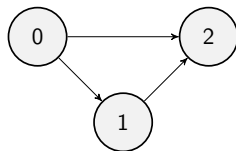
- the set of **TRs** is a set of nodes $R_G = N$
- a path p **covers** a node n if p visits n

Definition (Edge Coverage (EC))

The **edge coverage** criterion $C_G = (R_G, \sim)$ is defined with:

- the set of **TRs** is a set of edges $R_G = E$
- a path p **covers** an edge (n, m) if p visits (n, m)

NC and EC are only different when there is an edge and another subpath between a pair of nodes (e.g., an **if-else** statement).



Definition (k -Limiting Path Coverage (k -PC))

The k -**limiting path coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of paths whose lengths are bounded by k :

$$R_G = \{p \in P_G \mid |p| \leq k\}$$

- a path p **covers** a path q if q is a subpath of p :

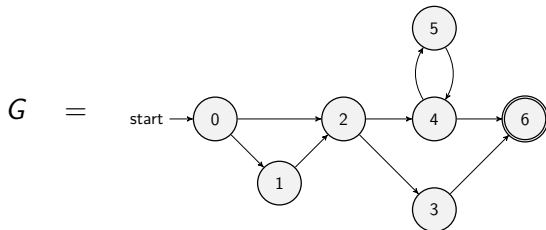
$$p \sim q \iff q \preceq p$$

Definition (Edge-Pair Coverage (EPC))

The **edge-pair coverage** criterion is 2-limiting path coverage.

Definition (Complete Path Coverage (CPC))

The **complete path coverage** criterion is ∞ -limiting path coverage.



- Node Coverage (NC)

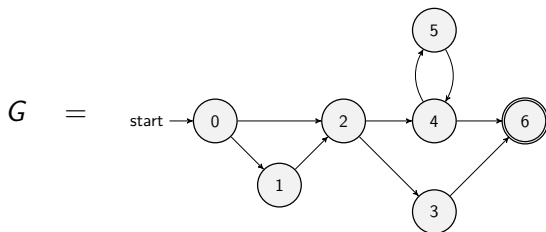
TRs $R_G = \{0, 1, 2, 3, 4, 5, 6\}$

Test Paths = $\{[0, 1, 2, 3, 6], [0, 1, 2, 4, 5, 4, 6]\}$

- Edge Coverage (EC)

TRs $R_G = \{\dots, (0, 1), (0, 2), (1, 2), (2, 3), (2, 4), (3, 6), (4, 5), (4, 6)\}$

Test Paths = $\{[0, 1, 2, 3, 6], [0, 1, 2, 4, 5, 4, 6]\}$



- Edge-Pair Coverage (EPC)

TRs $R_G = \{$
 $\dots, [0, 1, 2], [0, 2, 3], [0, 2, 4], [1, 2, 3], [1, 2, 4], [2, 3, 6],$
 $[2, 4, 5], [2, 4, 6], [4, 5, 4], [5, 4, 5], [5, 4, 6]$
 $\}$

Test Paths = $\{$
 $[0, 1, 2, 3, 6], [0, 1, 2, 4, 6], [0, 2, 3, 6], [0, 2, 4, 5, 4, 5, 4, 6]$
 $\}$

- If a graph contains a loop, it has an **infinite** number of paths
- In this case, the **complete path coverage** (CPC) is not feasible.
- One possible solution is to use a **specified path coverage (SPC)** criterion with a set of paths manually specified by the tester.
- However, it is highly **dependent** on the tester's **expertise**.
- Attempts to deal with loops:
 - 1970s: Execute cycles once ([4, 5, 4] in the previous example)
 - 1980s: Execute each loop, exactly once
 - 1990s: Execute loops 0 times, once, more than once
 - 2000s: **Prime paths**

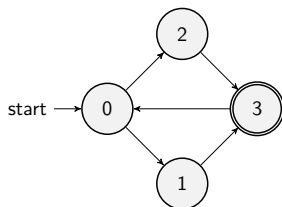
Definition (Simple Path)

A **simple path** is a path that does not contain a repeated node, except for the start and final nodes. In other words,

- No internal loops
- A loop is a simple path

Definition (Prime Path)

A **prime path** is a simple path that is not a subpath of other simple paths.

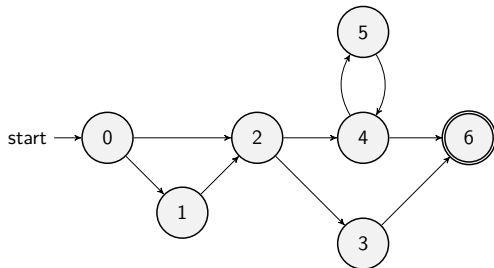


- **23 Simple Paths:**

[0], [1], [2], [3], [0, 1], [0, 2], [1, 3], [2, 3], [3, 0],
 [0, 1, 3], [0, 2, 3], [1, 3, 0], [2, 3, 0], [3, 0, 1], [3, 0, 2],
 [0, 1, 3, 0], [0, 2, 3, 0], [1, 3, 0, 1], [1, 3, 0, 2],
 [2, 3, 0, 1], [2, 3, 0, 2], [3, 0, 1, 3], [3, 0, 2, 3]

- **8 Prime Paths:**

[0, 1, 3, 0], [0, 2, 3, 0], [1, 3, 0, 1], [1, 3, 0, 2],
 [2, 3, 0, 1], [2, 3, 0, 2], [3, 0, 1, 3], [3, 0, 2, 3]



- **38 Simple Paths**
- **9 Prime Paths:**

[0, 1, 2, 3, 6] [5, 4, 6]

[0, 2, 4, 6] executes the loop **0** times

[0, 1, 2, 4, 5] [4, 5, 4]

[4, 5, 4] executes the loop **once**

[0, 1, 2, 4, 6] [5, 4, 5]

[5, 4, 5] executes the loop **more than once**

[0, 2, 3, 6]

[0, 2, 4, 5]

[0, 2, 4, 6]

Definition (Prime Path Coverage (PPC))

The **prime path coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of prime paths:

$$R_G = \{p \in P_G \mid p \text{ is a prime path}\}$$

- a path p **covers** a prime path q if q is a subpath of p :

$$p \sim q \iff q \preceq p$$

Definition (Round-Trip Path)

A **round-trip path** is a prime path that starts and ends at the same node.

Definition (Complete Round-Trip Path Coverage (CRPC))

The **complete round-trip path coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of all round-trip paths:

$$R_G = \{p \in P_G \mid p \text{ is a round-trip path}\}$$

- a path p **covers** a round-trip path q if q is a subpath of p :

$$p \sim q \iff q \preceq p$$

Definition (Simple Round-Trip Path Coverage (SRPC))

The **simple round-trip path coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of nodes visited by at least one round-trip path:

$$R_G = \{n \in N \mid \exists p \in P_G. p \text{ is a round-trip path} \wedge n \in p\}$$

- a path p **covers** a node n if at least one round-trip path for n is a subpath of p :

$$p \sim n \iff \exists q \in P_G. q \text{ is a round-trip path} \wedge n \in q \wedge q \preceq p$$

- CRPC and SRPC omit nodes and edges not in round-trip paths
- In other words, they only focus on loops

Prime paths do not have **internal loops!**

Definition (Tour)

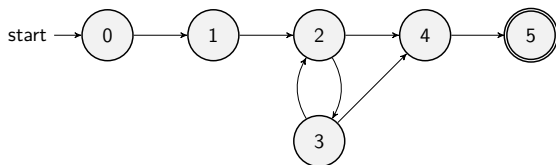
A test pest p **tours** a path q if q is a subpath of p

Definition (Tour with Sidetrips)

A test pest p **tours** a path q with **sidetrips** if and only if every **edge** in q is also in p in **the same order**.

Definition (Tour with Detours)

A test pest p **tours** a path q with **detours** if and only if every **node** in q is also in p in **the same order**.



$[0, 1, 2, 4, 5]$ **tours** $[1, 2, 4]$

$[0, 1, 2, 3, 2, 4, 5]$ does **not tour** $[1, 2, 4]$ but **tours** it with **sidetrips** ($[2, 3, 2]$ is a sidetrip)

$[0, 1, 2, 3, 4, 5]$ does **not tour** $[1, 2, 4]$ with sidetrips but **tours** it with **detours** ($[2, 3, 4]$ is a detour)

- An **infeasible** test requirement **cannot be satisfied**
 - Unreachable statements (dead code)
 - A subpath that can only be executed with a contradiction
- Most **coverage criteria** have some infeasible test requirements
- It is usually **undecidable** whether all test requirements are feasible
- When **sidetrips** or **detours** are not allowed, many structural coverage criteria have **more infeasible test requirements**
- **Practical solutions:** (1) try to satisfy as many test requirements as possible **without** sidetrips or detours, (2) **allow** sidetrips or detours to try to satisfy not yet satisfied test requirements

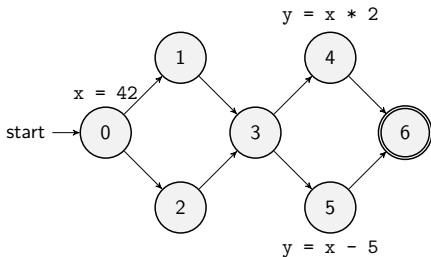
Our goal is to **ensure** that values are computed and **used** correctly.

Definition (Def)

A **definition** of a variable is a **location** in the program where a value is assigned to the variable.

Definition (Use)

A **use** of a variable is a **location** in the program where the value of the variable is accessed.



$$def(0) = \{x\}$$

$$def(4) = \{y\}$$

$$def(5) = \{y\}$$

$$use(4) = \{x\}$$

$$use(5) = \{x\}$$

Definition (DU-Pair)

A **du-pair** is a pair of a locations (l, l') such that a variable x is defined at l and used at l' .

Definition (Def-Clear Path)

A path from l to l' is **def-clear** with respect to a variable x if x is not given another value on any of the nodes or edges in the path.

Definition (DU-Path)

A **du-path** is a simple subpath that is def-clear with respect to x from a def of x to a use of x .

- $du(n, n', x)$ is the set of du-paths from n to n' with respect to x
- $du(n, x)$ is the set of du-paths from n to any use of x

Definition (DU-Tour)

A test path p **du-tours** a du-path q with respect to x if p tours d and the subpath taken is def-clear with respect to x

Sidetrips or **detours** can be used, just as with previous touring.

Definition (All-Defs Coverage (ADC))

The **all-defs coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of pairs of nodes and variables such that

$$R_G = \{(n, x) \mid n \in N \wedge |du(n, x)| > 0\}$$

- a path p **covers** a pair (n, x) if p du-tours a du-path in $du(n, x)$

$$p \sim (n, x) \iff \exists q \in du(n, x). p \text{ du-tours } q$$

Definition (All-Uses Coverage (AUC))

The **all-uses coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of triples of two nodes and variables such that

$$R_G = \{(n, n', x) \mid n, n' \in N \wedge |du(n, n', x)| > 0\}$$

- a path p **covers** (n, n', x) if p du-tours a du-path in $du(n, n', x)$

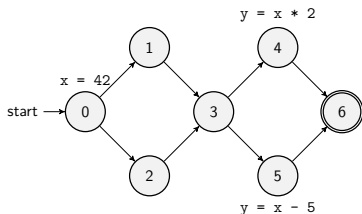
$$p \sim (n, n', x) \iff \exists q \in du(n, n', x). p \text{ du-tours } q$$

Definition (All-DU-Paths Coverage (ADUPC))

The **all-du-paths coverage** criterion $C_G = (R_G, \sim)$ is

- the set of **TRs** is a set of du-paths: $R_G = \{q \in P_G \mid q \text{ is a du-path}\}$
- a path p **covers** a du-path q if p du-tours q :

$$p \sim q \iff p \text{ du-tours } q$$



- **All-Defs Coverage (ADC)**

$$\text{TRs } R_G = \{(0, x)\}$$

$$\text{Test Paths} = \{[0, 1, 3, 4, 6]\}$$

- **All-Uses Coverage (AUC)**

$$\text{TRs } R_G = \{(0, 4, x), (0, 5, x)\}$$

$$\text{Test Paths} = \{[0, 1, 3, 4, 6], [0, 1, 3, 5, 6]\}$$

- **All-DU-Paths Coverage (ADUPC)**

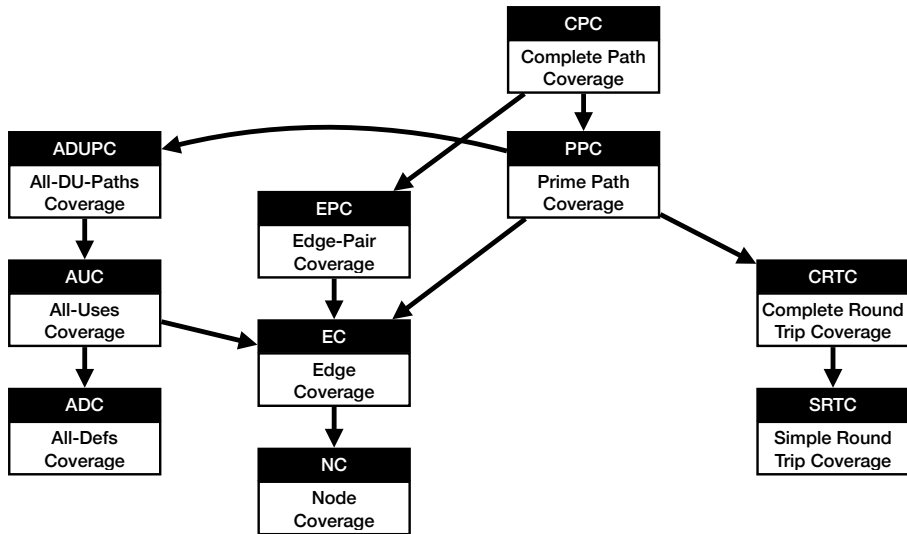
$$\text{TRs } R_G = \{[0, 1, 3, 4], [0, 1, 3, 5], [0, 2, 3, 4], [0, 2, 3, 5]\}$$

$$\text{Test Paths} = \{[0, 1, 3, 4, 6], [0, 1, 3, 5, 6], [0, 2, 3, 4, 6], [0, 2, 3, 5, 6]\}$$

Definition (Subsumption)

A coverage criterion $C_G = (R_G, \sim)$ **subsumes** another coverage criterion $C'_G = (R'_G, \sim')$ if and only if any test suite T satisfying C_G satisfies C'_G .

- **Edge Coverage (EC)** subsumes **Node Coverage (NC)**
- **Edge-Pair Coverage (EPC)** subsumes **Edge Coverage (EC)**
- **Prime Path Coverage (PPC)** subsumes **Edge Coverage (EC)**
- However, **Prime Path Coverage (PPC)** does **not** subsume **Edge-Pair Coverage (EPC)**
 - If a node n has a self-cycle, EPC will require non-prime path $[n, n, m]$
- **Complete Round-Trip Path Coverage (CRPC)** and **Simple Round-Trip Path Coverage (SRPC)** do not subsume **Node Coverage (NC)**



1. Graph Coverage

Structural Coverage

Data-Flow Coverage

Subsumption Relationships

2. Logic Coverage

Simple Logic Expression Coverage

Active Clause Coverage

Inactive Clause Coverage

Subsumption Relationships

3. Neuron Coverage

4. Feature-Sensitive Coverage

- Logic expressions show up in many situations
- Covering logic expressions is required by US Federal Aviation Administration for safety critical software
- Logical expressions can come from many sources
- Tests are intended to choose some subset of the total number of truth assignments to the expressions

- A **predicate** is an expression that evaluates to a **boolean** value
- Predicates can contain
 - **Boolean variables**
 - non-boolean variables with a **comparison** operator
 - Boolean **function** calls
- Internal structure is created by logical operators
 - \neg – the **negation** operator
 - \wedge – the **conjunction** operator
 - \vee – the **disjunction** operator
 - \oplus – the **exclusive or** operator
 - \Rightarrow – the **implication** operator
 - \Leftrightarrow – the **equivalence** operator
- A **clause** is a predicate without logical operators

$$(a < b) \vee f(z) \wedge D \wedge (m \geq n \times o)$$

- Four clauses:
 - $(a < b)$ – relational expression
 - $f(z)$ – boolean-value function
 - D – boolean variable
 - $(m \geq n \times o)$ – relational expression
- Most predicates have few clauses

Abbreviations:

- P is the set of **predicates**
- p is a single **predicate** in P
- C is the set of **clauses**
- C_p is the set of **clauses** in predicate p
- c is a single **clause** in C

Definition (Predicate Coverage (PC))

For each predicate $p \in P$, test requirements are the **truth** or **falsity** of p .

It is sometimes called **decision coverage**.

Definition (Clause Coverage (CC))

For each clause $c \in C$, test requirements are the **truth** or **falsity** of c .

It is sometimes called **condition coverage**.

$$p = ((a < b) \vee D) \wedge (m \geq n \times o)$$

- Predicate Coverage (PC)**

<i>a</i>	<i>b</i>	<i>D</i>	<i>m</i>	<i>n</i>	<i>o</i>	<i>p</i>
5	10	<i>T</i>	1	1	1	<i>T</i>
10	5	<i>F</i>	1	1	1	<i>F</i>

- Clause Coverage (CC)**

<i>a</i>	<i>b</i>	<i>D</i>	<i>m</i>	<i>n</i>	<i>o</i>	$(a < b)$	<i>D</i>	$(m \geq n \times o)$
5	10	<i>F</i>	1	1	1	<i>T</i>	<i>F</i>	<i>T</i>
10	5	<i>T</i>	1	2	2	<i>F</i>	<i>T</i>	<i>F</i>

- PC does not fully exercise all the clauses in the predicate, especially in the presence of **short circuit** evaluation
- CC does not subsume PC. For example, the following test suite satisfies CC but not PC:

	A	B	$A \vee B$
$A \vee B$	T	F	T
	F	T	T

- The simplest solution is to test **all combinations**

- PC does not fully exercise all the clauses in the predicate, especially in the presence of **short circuit** evaluation
- CC does not subsume PC. For example, the following test suite satisfies CC but not PC:

$$A \vee B$$

A	B	$A \vee B$
T	F	T
F	T	T

- The simplest solution is to test **all combinations**

Definition (Combinatorial Coverage (CoC))

For each predicate $p \in P$, test requirements in **combinatorial coverage (CoC)** are the all **combinations** of the truth or falsity of the clauses in C_p .

For example, we need the following combinations for the predicate p

$$p = ((a < b) \vee D) \wedge (m \geq n \times o)$$

$(a < b)$	D	$(m \geq n \times o)$	p
T	T	T	T
T	T	F	F
T	F	T	T
T	F	F	F
F	T	T	T
F	T	F	F
F	F	T	F
F	F	F	F

- This is simple, neat, clean, but quite **expensive!**
- 2^N tests required for N clauses
- Let's test each clause **independently** from the other clauses using whether each clause **determines** the entire predicate.

Definition (Determination)

A clause c in predicate p , called the **major clause**, **determines** p if and only if the values of the remaining **minor clauses** c' are such that changing the value of p .

- A (or B) **determines** $A \vee B$ if B (or A) is **false**, and A (or B) **determines** $A \wedge B$ if B (or A) is **true**.

A	B	$A \vee B$
T	T	T
T	F	T
F	T	T
F	F	F

A	B	$A \wedge B$
T	T	T
T	F	F
F	T	F
F	F	F

- A (or B) always **determines** $A \oplus B$, and A (or B) always **determines** $A \Leftrightarrow B$.

A	B	$A \oplus B$
T	T	F
T	F	T
F	T	T
F	F	F

A	B	$A \Leftrightarrow B$
T	T	T
T	F	F
F	T	F
F	F	T

Definition (Active Clause Coverage (ACC))

For each predicate $p \in P$, test requirements in **active clause coverage (ACC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ determine p and (2) the truth or falsity of c .

- For example, $p = A \vee B$, and pick A (or B) as the major clause.
 - ① $A = \text{true}$ and $B = \text{false}$ (A determines p)
 - ② $A = \text{false}$ and $B = \text{false}$ (A determines p)
 - ③ $A = \text{false}$ and $B = \text{true}$ (B determines p)
 - ④ ~~$A = \text{false}$ and $B = \text{false}$ (B determines p)~~
- The last one is **duplicate** and can be omitted.
- Another name of ACC is **modified condition/decision coverage (MC/DC)**, which is required by the **US Federal Aviation Administration** for **safety critical software**.

- **Ambiguity** – Do the **minor clauses** have to have the **same values** when the **major clause** is **true** or **false**?
- For example, $p = A \vee (B \wedge C)$, and pick A as the major clause.
 - ① $A = \text{true}$, $B = \text{false}$, $C = \text{true}$
 - ② $A = \text{false}$, $B = \text{false}$, $C = \text{false}$ (is $C = \text{false}$ allowed?)
- This question caused a **confusion** among testers for years
- Consider this carefully leads to three separate coverage criteria:
 - Minor clauses **do not need** to be the same
 - Minor clauses **must** be the same
 - Minor clauses **force the predicate** to have different values

Definition (General Active Clause Coverage (GACC))

For each predicate $p \in P$, test requirements in **general active clause coverage (GACC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ determine p and (2) the truth or falsity of p . The values chosen for the **minor clauses** do **not need** to be the same when the **major clause** is **true** or **false**.

- Unfortunately, GACC does **not subsume predicate coverage (PC)**.
- For example, the following selection satisfies GACC but not PC:

A	B	$A \leftrightarrow B$
T	T	T
T	F	F
F	T	F
F	F	T

Definition (Restricted Active Clause Coverage (RACC))

For each predicate $p \in P$, test requirements in **restricted active clause coverage (RACC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ determine p and (2) the truth or falsity of p . The values chosen for the **minor clauses** **must** be the same when the **major clause** is **true** or **false**.

- This has been a **common interpretation** by aviation developers
- RACC often leads to **infeasible test requirements**
- There is **no logical reason** for such a strict restriction
- Our **goal** is just to subsume **predicate coverage (PC)**

Definition (Correlated Active Clause Coverage (CACC))

For each predicate $p \in P$, test requirements in **correlated active clause coverage (CACC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ determine p and (2) the truth or falsity of p . The values chosen for the **minor clauses** **force the predicate** to have different values when the **major clause** is **true** or **false**.

- A **more recent** interpretation
- CACC **implicitly** allows minor clauses to have different values
- CACC **explicitly** subsumes **predicate coverage (PC)**

A	B	C	$A \wedge (B \vee C)$
T	T	T	T
F	T	T	F
T	T	F	T
F	T	F	F
T	F	T	T
F	F	T	F

A	B	C	$A \wedge (B \vee C)$
T	T	T	T
T	T	F	T
T	F	T	T
F	T	T	F
F	T	F	F
F	F	T	F

- We pick A as the major clause
- The left table shows that there are only **three combinations** allowed in **RACC**
- The right table shows that there are **nine combinations** allowed in **CACC** by selecting any three cases for each truth value of A

- **Inactive clause coverage (ICC)** is the **dual** of **active clause coverage (ACC)**
- ACC criteria ensure that **major** clauses **determine** the predicate
- ICC criteria ensure that **major** clauses **do not determine** the predicate

Definition (Inactive Clause Coverage (ICC))

For each predicate $p \in P$, test requirements in **inactive clause coverage (ICC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ **not determine** p and (2) the truth or falsity of c .

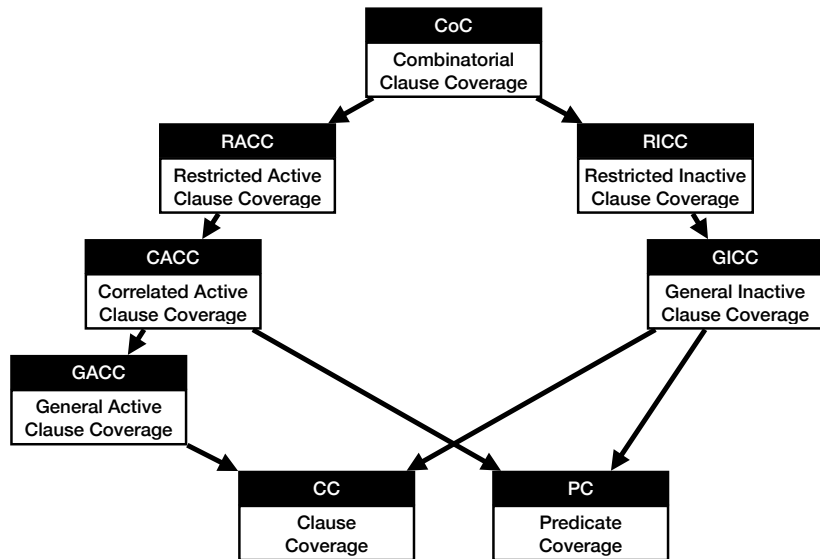
Definition (General Inactive Clause Coverage (GICC))

For each predicate $p \in P$, test requirements in **general inactive clause coverage (GICC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ **not determine** p and (2) the truth or falsity of c . The values chosen for the **minor clauses** do **not need** to be the same when the **major clause** is **true** or **false**.

Definition (Restricted Inactive Clause Coverage (RICC))

For each predicate $p \in P$, test requirements in **restricted inactive clause coverage (RICC)** are pairs of (1) **conditions** that make each selected **major clause** $c \in C_p$ **not determine** p and (2) the truth or falsity of c . The values chosen for the **minor clauses** **must** be the same when the **major clause** is **true** or **false**.

- Unlike ACC, the notion of **correlation** is not relevant to ICC (major clause c does not determine p , so cannot correlate with it)



1. Graph Coverage

Structural Coverage

Data-Flow Coverage

Subsumption Relationships

2. Logic Coverage

Simple Logic Expression Coverage

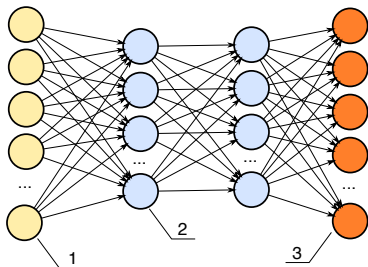
Active Clause Coverage

Inactive Clause Coverage

Subsumption Relationships

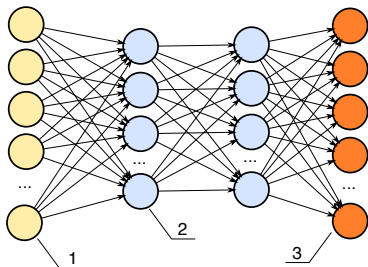
3. Neuron Coverage

4. Feature-Sensitive Coverage



How to define coverage criteria for **deep neural networks (DNN)**?

Neuron Coverage!



- A **DNN classifier** can be formalized as a function $f : X \rightarrow Y$, a mapping from a set of inputs X into a set of labels Y .
- The output of the DNN classifier is a **probability distribution** $P(Y | x)$, which is the probability that an input vector $x \in X$ belongs to each class of labels in Y .
- A DNN classifier f usually contains an input layer, a number of hidden layers, and an output layer; each layer consists of many **neurons**.

- The **parameters** θ of the DNN are the **weights** of each connected edge between neurons of two adjacent layers.
- For an input vector x , the output of DNN $f_{\theta}(x)$ can be computed as the **weighted sum** of the outputs of all the neurons.
- Given a training dataset $D = \{(x_i, y_i)\}_{i=1}^N$, the goal is to learn to optimize the following **objective**:

$$\min_{\theta} \frac{1}{N} \sum_{i=1}^N \mathcal{L}(f_{\theta}(x_i), y_i)$$

where \mathcal{L} is a **loss function** to calculate the penalties for incorrect classifications.

- For a neuron n in a DNN model f , the **output** of n for an input x is denoted as $f_{\theta}(n, x)$.

Definition (Activate Neuron)

A neuron is **activated** if the weighted sum of its inputs exceeds a certain threshold t .

$$AC(x, t) = \{n \mid f_{\theta}(n, x) > t\}$$

Definition (Neuron Coverage (NC))

For a DNN model f and a threshold t , the test requirements in **neuron coverage (NC)** are the **activation** of each neuron n in the DNN model f .

$$NC(T, t) = \frac{|\{n \mid \exists x \in T. f_{\theta}(n, x) > t\}|}{|N|}$$

For a neuron n , the lower and upper boundary of its output values on **training data** can be denoted as low_n and up_n , respectively.

Definition (k -Multisection Neuron Coverage (KMNC))

For a DNN model f and a number of sections k , the test requirements in **k -multisection neuron coverage (KMNC)** are k **equal sections** of $[low_n, up_n]$ for each neuron n in the DNN model f .

$$KMNC = \frac{\sum_{n \in N} |\{S_m^n \mid \exists x \in T. f_\theta(n, x) \in S_m^n\}|}{|N|}$$

where

$$S_m^n = \left[low_n + \frac{m \times (up_n - low_n)}{k}, low_n + \frac{(m + 1) \times (up_n - low_n)}{k} \right]$$

For new test inputs T , the output values of neurons may fall into $(-\infty, low_n)$ or $(up_n, +\infty)$. instead of the derived boundary $[low_n, up_n]$.

Definition (Upper or Lower Neuron Coverage (UNC or LNC))

For a DNN model f and a number of sections k , the test requirements in **upper neuron coverage (UNC)** (or **lower neuron coverage (LNC)**) are the **upper boundary** (or **lower boundary**) of the output values of each neuron n in the DNN model f .

$$UNC = \frac{|\{n \mid \exists x \in T. f_{\theta}(n, x) > up_n\}|}{|N|}$$

$$LNC = \frac{|\{n \mid \exists x \in T. f_{\theta}(n, x) < low_n\}|}{|N|}$$

UNC is often called **strong neuron activation coverage (SNAC)**.
Neuron Boundary Coverage (NBC) is combination of UNC and LNC.

$$NBC = (|UNC| + |LNC|)/2$$

Definition (Top- k Neuron Coverage (TKNC))

TKNC is a layer-level coverage testing criterion that measures the ratio of neurons that have at least been the **most active k neurons** of **each layer** on a given test set T once.

$$TKNC = \frac{|\bigcup_{x \in T} \bigcup_{1 \leq l \leq L} top_k(x, l)|}{|N|}$$

where L denotes the number of layers in the DNN model f , and $top_k(x, l)$ denotes the neurons which have largest k output values in the l -th layer of the DNN model f for the input x .

We can extend this criterion to the **Top- k Neuron Pattern Coverage (TKNPC)** by considering the **combination** of the most active neurons in each layer.

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- **[PLDI'23]** J. Park et al. *“Feature-Sensitive Coverage for Conformance Testing of Programming Language Implementations.”*
- It suggests a new way to refine a given graph coverage criterion using **feature-sensitive** information.
- Slides: [link](#)

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- Coverage Criteria (Homework)

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